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The Compact Euclidean Space Forms
of Dimension Four

By

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**The Compact Euclidean Space Forms
of Dimension Four**

Ronald Dorian Levine

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of Dimension Four

Abstract

Ronald Dorian Levine

The main result of this work is a complete affine classification of the compact four-dimensional riemannian manifolds of constant zero curvature. There are 75 classes of which 27 are orientable. These spaces are described by means of tables of their algebraic invariants: fundamental groups, first homology groups, and linear holonomy groups. A second feature of this work is the description of a set of computer programs which were instrumental in obtaining the first result and which promise to be useful in extending the work to higher dimensions and in the more general problem of enumerating the higher dimensional crystallographic groups.

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Chapter 0

Introduction

The main result of this work is a complete affine classification of the 4-dimensional compact riemannian manifolds of constant zero curvature.^(*) There are 75 classes, of which 27 are orientable. Tables C and D of the appendix contain their descriptions by means of their algebraic invariants: fundamental groups, first homology groups, and linear holonomy groups. A second feature of this paper is the description of a set of computer programs which were instrumental in obtaining the first result and which promise to be useful in extending the work to higher dimensions and in attacking the more general problem of enumerating higher dimensional crystallographic groups.

The structure theory of riemannian manifolds of zero curvature, also called flat manifolds or euclidean space forms, as well as the history of its development, may be found in the book by Wolf. [24] ^(**) Here we summarize the principal results for the compact euclidean space forms.

(*) In this paper the word "manifold" is understood to entail "differentiable" and "connected".

(**) Digits in square brackets refer to the bibliography on p. 44.

The most satisfactory equivalence relation for classifying flat manifolds is affine equivalence, i.e. equivalence by diffeomorphisms preserving the affine connection, but in the compact case it turns out that this coincides with classification by homeomorphism. Every n -dimensional compact euclidean space form M can be realized as the quotient of euclidean R^n by the action of a discrete group G of isometries, with each nontrivial element of G having no fixed points in R^n . Then G , as the group of deck transformations of the universal covering of M , is isomorphic to its fundamental group $\pi_1(M)$. Conversely, if G is a discrete group of fixed-point-free isometries of R^n , then the quotient space R^n/G has naturally the structure of a flat riemannian manifold which is compact if and only if G contains a set of n linearly independent translations.

Moreover, two compact euclidean space forms are affinely equivalent if and only if their fundamental groups are isomorphic. For each dimension n , there are only finitely many classes of compact flat manifolds; in dimensions 1, 2, 3, there are 1, 2, 10 (resp.) of them, of which 1, 1, 6, (resp.) are orientable.

Discrete groups of isometries of R^n acting with compact quotient are called n -dimensional crystallographic groups, or space groups. In dimensions 3 and 2 they have famous applications to the description of physical crystal lattices and ornamental designs. It can be seen that, for these groups, acting without fixed points is equivalent to being torsion-free,

i.e., to having no non-trivial elements of finite order.

Thus, the problem of determining the compact euclidean space forms is reduced to determining all torsion-free crystallographic groups up to isomorphism.

The structure theory of crystallographic groups, due mainly to Bieberbach, can be found in [2, 3, 26]. Again we summarize the well-known main results on which the present computation is based.

Let G be an n -dimensional crystallographic group and let T be the set of translations in G . Then T is a normal subgroup of finite index in G , which is free abelian of rank n and maximal abelian in G . The finite quotient group $F = G/T$ is called the point group or crystal class of G , and in the torsion-free case F is isomorphic to the linear holonomy group of the corresponding flat manifold. Now conjugation of T by coset representatives in G of the elements of F induces a well-defined faithful representation of F in $\text{Aut}(T)$. The choice of an integral basis in T embeds T in \mathbb{Z}^n , and consequently F in $\text{GL}(n, \mathbb{Z})$. These embeddings can be extended to an embedding of G in $\text{GL}(n, \mathbb{Z}) \cdot \mathbb{Q}^n$ (semi-direct product), and for the remainder of this paper we will work with these groups in this concrete arithmetical form. Specifically, \mathbb{Q}^n denotes the additive group of n -component rational column vectors, \mathbb{Z}^n denotes the integer lattice in \mathbb{Q}^n , and $\text{GL}(n, \mathbb{Z})$ denotes the multiplicative group of $n \times n$ unimodular integer matrices. The semi-direct product of $\text{GL}(n, \mathbb{Z}) \cdot \mathbb{Q}^n$ consists of all pairs $(A, s) \in \text{GL}(n, \mathbb{Z}) \times \mathbb{Q}^n$, with the multiplication rule,

$$(A,s)(B,t) = (AB, s + At).$$

We will call A the linear part and s the translation part of (A,s) .

According to a theorem of Bieberbach, every isomorphism between two crystallographic groups is induced by an inner automorphism of $GL(n,Z) \cdot Q^n$, and this of course entails that the point groups be conjugate in $GL(n,Z)$. A class of finite subgroups of $GL(n,Z)$ conjugate in $GL(n,Z)$ is called an arithmetical crystal class, while a class of finite subgroups of $GL(n,Z)$ conjugate in $GL(n,Q)$ is called a geometrical crystal class. For any finite subgroup F of $GL(n,Z)$, we will call G a crystallographic extension of F if G is a subgroup of $GL(n,Z) \cdot Q^n$ and there is an exact sequence

$$0 \rightarrow Z^n \xrightarrow{i} G \xrightarrow{j} F \rightarrow 1$$

in which i is the canonical inclusion and j is the restriction to G of the canonical projection $GL(n,Z) \cdot Q^n \rightarrow GL(n,Z)$.

Our determination of the 4-dimensional compact euclidean space forms follows, in three parts out of four, a general program enunciated by Charlap in his study [6] of the problem of classifying, for each given abstract finite group ϕ , the flat manifolds with holonomy isomorphic to ϕ . The first step in Charlap's program, "determine all integral representations of ϕ ", is difficult and has been completed only for the most simple groups ϕ . [22] The integral representations of the cyclic groups of prime order have been classified by Reiner [21] in terms of the ideal class groups of the prime cyclotomic fields,

and Charlap has used this to classify the flat manifolds of all dimensions whose holonomies are cyclic of prime order.

In the present work, we replace Charlap's first step with

Step I: Find a set of mutually arithmetically inequivalent finite subgroups of $GL(4, \mathbb{Z})$ which contains all possible holonomy groups.

Specifically, as we will see in Chapter 3, any matrix group F which has a torsion-free crystallographic extension necessarily satisfies

Condition A: Every element of F has non-zero fixed space.

In Chapter 1 we show how to determine a complete set of arithmetical crystal classes satisfying Condition A. In Chapter 2 we show how to determine, for each F found in Step I, a complete set of crystallographic extensions of F , which are perhaps not mutually non-isomorphic. Namely, we have

Step II: Compute the second cohomology group $H^2(F, \mathbb{Z}^n)$.

In Chapter 3 we treat

Step III: Recognize the torsion-free groups arising in Step II.

In Chapter 4 we carry out

Step IV: Find a mutually non-isomorphic set from the groups selected in Step III.

In each of these chapters the corresponding step is reduced to the level of arithmetical computations which can be completed electro-mechanically and we name the computer codes which have been developed to perform them. Further details on the codes

and the mechanics of their use are collected in Chapter 5, and listings of the programs comprise Table E.

Our central computational tool is the Fundamental Theorem of Finitely Generated Abelian Groups, which may be stated in the arithmetical form: If M is any $p \times q$ matrix with integer coefficients, then there exists a $p \times p$ unimodular integer matrix L , and a $q \times q$ unimodular integer matrix K such that

$$LMK = \begin{pmatrix} e_1 & & & & 0 \\ & e_2 & & & \\ & & \dots & & \\ & & & e_r & \\ 0 & & & & \dots & 0 \end{pmatrix}$$

and the integers r, e_1, \dots, e_r are uniquely determined by

$$0 < e_1 \mid e_2 \mid \dots \mid e_r .$$

The standard constructive proof may be found in [13]. We use this construction at every turn, and will refer to it in the text by the acronym FTAG. One version of our Compass language code for it, called DIAR, is listed in Table E.

In 1957, Calabi [5] announced a determination of the 4-dimensional compact euclidean space forms using a recursive geometric approach which is quite different from our method and which is described in [24]. He gives details neither of the computation, nor of the results, but only the numbers of orientable and non-orientable spaces occurring, and these numbers are contradicted by the present work. Charlap and Sah have privately circulated a report [7] of another determination made in the spirit of Charlap's program, but without the aid of a computer.

They give no details of the computations, but do give a final list of groups, which we find to be incorrect in several respects.

The author wishes to express his thanks to Professor Joseph A. Wolf who introduced him to the problem, and to Mr. D. J. Underwood and the staff of the Humboldt State College Computer Center who arranged his easy access to their facilities.

Chapter 1

The Arithmetical Classes of Holonomy

The crystal classes in dimensions $n < 4$ have been known for some time and are given in Burckhardt's book [3]. In dimension 2 there are 10 geometrical and 13 arithmetical classes, while in dimension 3 there are 32 geometrical and 73 arithmetical classes.

The results in dimension 4 begin in 1889 with Goursat [10], who gave a determination of the (infinitely many) conjugate classes of finite subgroups of $SO(4)$ and of those of $O(4)$ containing the inversion in the origin $-I$. This work was later corrected and extended by Seifert and Threlfall [23]. In these works the principal tool used for pulling up the 3-dimensional groups is the double covering $SO(4) \rightarrow SO(3) \times SO(3)$. In 1951, Hurley [11] extracted from Goursat's list of classes all those having integral representations, i.e. geometrical crystal classes of R^4 , and found the classes not containing $-I$ as subgroups of those containing $-I$. As amended in [12], there are 227 4-dimensional geometrical crystal classes.

The nine maximal 4-dimensional arithmetical crystal classes are worked out in a remarkable paper by Dade [9]. Finally, Neubüser and Bülow [4] have reported a computer search for the subgroups of Dade's groups, finding that there are 720 4-dimensional arithmetical crystal classes in all. However these authors have not yet released their list of groups.

We begin the present computation by extracting from Hurley's list of geometrical classes a list of all those which satisfy Condition A. Since Hurley names a finite matrix group by enumerating its elements by similarity type, this task is quite easy. There are 45 geometrical classes satisfying Condition A, of which 11 are orientation preserving. They are listed in Table A. Now we need to determine for each of these 45 classes the classes into which it splits under arithmetical equivalence, and this task is facilitated by the fact that every one of the 45 classes is $(3+1)$ -reducible.

We proceed according to the theory of [25] and [8] to build the 4-dimensional arithmetical crystal classes from the known 3-dimensional classification. Here we inject some of the language of representation theory, and it is well to note that the usual sense of equivalence for integral representations is not the same as our arithmetical equivalence. Namely, two integral representations $g, h: \phi \rightarrow GL(n, Z)$ of a group ϕ are equivalent in the sense of representation theory if there is an inner automorphism i_χ of $GL(n, Z)$ such that the diagram

$$\begin{array}{ccc} & g & \\ \phi & \rightarrow & GL(n, Z) \\ & & \uparrow i_\chi \\ & h & \\ \phi & \rightarrow & GL(n, Z) \end{array}$$

commutes, while for the arithmetical equivalence of $g(\phi)$ with $h(\phi)$ an arbitrary automorphism of ϕ can be substituted for the identity map in the diagram.

Now let F be one of our 45 candidates for holonomy groups.

Concretely we choose F to be a (3+1)-decomposed integral matrix group^(*) whose 3x3 block is one of Burckhardt's representatives of a 3-dimensional crystal class. We write

$$A = \begin{pmatrix} g(A) & 0 \\ 0 & h(A) \end{pmatrix}, \quad A \in F$$

where g is a 3-dimensional and h a 1-dimensional integral representation of F . Note that g, h are neither unique, nor, in general, faithful. Regarding $g(F)$ as a 3-dimensional geometrical crystal class, we get from Burckhardt a complete set of representatives g_1, g_2, \dots, g_r of the arithmetical classes into which it splits. Now it follows from the theory of [8] that every arithmetical class belonging to F can be represented in the form

$$(1) \quad A \mapsto \begin{pmatrix} g_i(A) & t(A) \\ 0 & h(A) \end{pmatrix} \quad \text{for some } i = 1, \dots, r$$

where $t: F \rightarrow \{3 \times 1 \text{ integer matrices}\}$ is a function which satisfies

$$(2) \quad \forall A, B \in F \quad t(AB) = g_i(A)t(B) + t(A)h(B)$$

and is called a binding function for g_i and h .

Moreover, for a given g_i and h , we get a finite, complete, but perhaps not mutually inequivalent set of representations of the form (1) by taking for binding functions a set of coset representatives of the cohomology group,

(*) The translation of Hurley's notation given in [18] was useful in constructing these.

$$(3) \quad C(g_i, h) = B(g_i, h) / B'(g_i, h)$$

where $B(g_i, h)$ is the additive group of all binding functions and $B'(g_i, h)$ is the subgroup of all inner binding functions, i.e. all those of the form

$$(4) \quad t(A) = g_i(A)d - dh(A) \quad \forall A \in F$$

where d is an arbitrary 3×1 integer matrix.

Computing this group for each g_i , $i = 1, \dots, r$, gives a list $\{F_1, \dots, F_s\}$ of matrix groups which contains all the arithmetical classes belonging to the geometrical class of F and which now must be tested for equivalences. For each pair (i, j) we must determine whether there exists a matrix $X \in GL(4, Z)$ with $F_i X = X F_j$. Note that equivalences may obtain between F_i arising from distinct g_j .

The computation of the cohomology group is made automatically and the determination of equivalence is made semi-automatically. We describe these computations in the context of binding together an arbitrary p -dimensional integral representation g of a finite group ϕ with an arbitrary q -dimensional integral representation h , for the programs are written in this generality.

We identify the additive group of all $p \times q$ integer matrices with Z^{pq} , and the set of all functions $\phi \rightarrow Z^{pq}$ with $Z^{k pq}$, where k is the order of ϕ . For each pair $(a, b) \in \phi \times \phi$ the binding function condition (2) gives a linear map

$$\begin{aligned} Z^{k pq} &\longrightarrow Z^{pq} \\ t &\longmapsto t(ab) - g(a)t(b) - t(a)h(b) \end{aligned}$$

and taking the direct sum over all pairs $(a, b) \in \phi \times \phi$ gives a linear map

$$\mathbb{Z}^{kpq} \xrightarrow{\epsilon} \mathbb{Z}^{k^2pq}$$

whose kernel is $B(g,h)$. Now sending an arbitrary pxq matrix d to the function (4) gives a linear map

$$(5) \quad \mathbb{Z}^{pq} \xrightarrow{\delta} \mathbb{Z}^{kpq}$$

and we have

$$C(g,h) = B(g,h)/B'(g,h) = \ker \epsilon / \text{im } \delta .$$

We claim that all information about $C(g,h)$ is contained in the map δ and the map ϵ can be ignored. For $B(g,h)$ is the kernel of a map into a free abelian group and therefore is a direct summand of \mathbb{Z}^{kpq} ; and since the theory guarantees that $C(g,h)$ is finite, $B'(g,h)$ spans $B(g,h)$ over Q (in fact, over $(1/k)\mathbb{Z}$). In other words

$$(6) \quad C(g,h) = ((Q \cdot \text{im } \delta) \wedge \mathbb{Z}^{kpq}) / \text{im } \delta .$$

Suppose ϕ is generated by $\{a_1, \dots, a_\mu\}$. Since a binding function is determined by its values on the generators, we can replace \mathbb{Z}^{kpq} in (5) and (6) with $\mathbb{Z}^{\mu pq}$. Now we represent δ by a $(\mu pq) \times (pq)$ integer matrix Δ , which is easily constructed from the input data $\{g(a_i), h(a_i), i = 1, \dots, \mu\}$ according to (4). The columns of Δ span $B'(g,h)$ over \mathbb{Z} and $B(g,h)$ over Q . Apply FTAG:

$$(7) \quad L \Delta K = \text{diag}(e_1, \dots, e_r, 0, \dots, 0).$$

Now it is clear from (6) and (7) that a complete set of representatives of the distinct elements of $C(g,h)$ is obtained by taking the elements (*)

$$L^{-1}[f_1, \dots, f_r, 0, \dots, 0] \quad 0 \leq f_i < e_i, \quad i = 1, \dots, r$$

(*) We reconcile the demands of typographical economy with our preference for writing operators on the left by adopting the convention that a row vector enclosed in square brackets is to be read as a column.

These vectors are then reassembled into sets of pxq matrices

$$t(a_i), \quad i = 1, \dots, \mu$$

and stuffed into the matrices

$$\begin{pmatrix} g(a_i) & * \\ 0 & h(a_i) \end{pmatrix} \quad i = 1, \dots, \mu,$$

to give the list of integral representations of ϕ which we seek.

Now let F_1 and F_2 be two integral representations of degree n . To settle the question of their arithmetical equivalence, we must determine whether for any automorphism σ of ϕ the equations

$$(8) \quad F_1(a_i)X - X F_2(\sigma(a_i)) = 0 \quad i = 1, \dots, \mu$$

have a solution by an nxn unimodular integer matrix X .

For each σ , (8) constitutes a set of μn^2 homogeneous linear equations in the n^2 coefficients of X . The solution set is a subgroup of the additive group of all nxn integer matrices and we find an integral basis for it by applying FTAG to the matrix of coefficients of (8). This allows us to express the general matrix satisfying (8) as a matrix Ξ of homogeneous linear forms in some number $s = n^2 - r$ of variables. (*)

Thereby, the problem of the equivalence of F_1 and F_2 is reduced to determining the existence of integer roots of one of the polynomial equations

$$(9) \quad \det \Xi = \pm 1$$

(*) See the discussion of equation (2a) of Chapter 4 for more details.

The left side of (9) is a homogeneous polynomial of degree n in s variables. No one knows a general mechanical procedure for determining the existence of integer roots of polynomials in several variables. Our computer program, in each case, computes the polynomial $\det E$ and then tests it for an obvious necessary condition, namely, that the greatest common divisor of the set of coefficients be 1. If this condition is satisfied, the polynomial is printed out as a problem for the user. In almost every such case in the present computation, the existence of integer roots was apparent by inspection; the few exceptions were not critical due to redundancy in the computations. (**)

Notice that not every abstract automorphism of ϕ must be tested in (8). First, if F_1 and F_2 are not equivalent as integral representations, i.e. with $\sigma = \text{id}$, but are equivalent by some other σ , then σ is necessarily an outer automorphism of ϕ . Second, in many specific cases some automorphisms of ϕ are seen to be forbidden by the numerical invariants of the matrices of F_1 and F_2 .

The computations just described are the function of the pair of programs VB/VC described in Chapter 5 and listed in Table E.

The 45 4-dimensional geometrical crystal classes of our initial sample are seen to split into 229 arithmetical classes. Of these, according to the subsequent computations, 49 are holonomy classes of compact euclidean space forms and are listed in Table C.

(**) The same semi-automatic method of deciding arithmetical equivalence was used in [4].

Recently, Janssen [14] has reported determining all the (3+1)-reducible arithmetical crystal classes, using a method rather more geometrical than arithmetical, constructing the Bravais lattices corresponding to these groups from the known 3-dimensional Bravais lattices. Differences in notation, representations, and viewpoint militate against an easy detailed comparison of our results with his; but a cursory comparison shows that we agree in the number of arithmetical classes into which each geometrical class splits.

Chapter 2

The Cohomology Group of Extensions

From now until the end of Chapter 4 we take F to be a fixed finite subgroup of $GL(n, Z)$. We would like to determine all crystallographic extensions G of F up to isomorphism. We begin in Step II by determining all possible G up to a classification which is finer than isomorphism. We call two extensions G and G_1 of F equivalent if there is a commuting diagram

$$\begin{array}{ccccccccc} 0 & \rightarrow & Z^n & \rightarrow & G & \rightarrow & F & \rightarrow & 0 \\ & & \uparrow & & \tau & & \uparrow & & \\ 0 & \rightarrow & Z^n & \rightarrow & G_1 & \rightarrow & F & \rightarrow & 0 \\ & & & & \downarrow & & & & \\ & & & & 1 & & & & \\ & & & & i & & & & \end{array}$$

and in this case τ is necessarily an isomorphism. (*)

As is well known [19], the set of all equivalence classes of extensions has naturally the structure of a finite abelian group, the second cohomology group of F in Z^n , denoted $H^2(F, Z^n)$. (**) Moreover, in our crystallographic context, the exact sequence of F -modules

$$0 \rightarrow Z^n \subset Q^n \xrightarrow{\pi} (Q/Z)^n \rightarrow 0,$$

(*) Our nomenclatures "isomorphic" and "equivalent" follow MacLane; Zassenhaus used "gewöhnlich äquivalent" and "stark äquivalent" for these, resp.; Burckhardt used "äquivalent" and "null-äquivalent".

(**) Actually true for extensions of more general type than we are considering.

the corresponding connecting homomorphism in cohomology, and the divisibility of Q^n (as an abelian group), combine to provide an isomorphism

$$H^2(F, Z^n) \approx H^1(F, (Q/Z)^n),$$

where the first cohomology group $H^1 = H^1(F, (Q/Z)^n)$ may be realized as a quotient $H^1 = Cr/Pr$, where $Cr = Cr(F, (Q/Z)^n)$ is the group of crossed homomorphisms,

$$Cr = \{F \xrightarrow{t} (Q/Z)^n : \forall A, B \in F, t(AB) = t(A) + At(B)\},$$

and $Pr = Pr(F, (Q/Z)^n)$ is the group of all principal crossed homomorphisms,

$$Pr = \{F \xrightarrow{t} (Q/Z)^n : \exists s \in (Q/Z)^n, \forall A \in F, t(A) = s - As\}$$

The crystallographic group determined by a crossed homomorphism t is the set of all elements of $GL(n, Z) \cdot Q^n$ of the form $(A, t(A) + z)$, $z \in Z^n$, $A \in F$. Here we have used the same symbol $t(A)$ for an element of $(Q/Z)^n$ and an arbitrary one of its lifts in Q^n .

Our computation of the cohomology group H^1 is similar to but more complicated than the cohomology calculation of the preceding chapter. The algorithm is due to Zassenhaus. [26] We begin by selecting a presentation for F :

$$\begin{aligned} \phi &= \langle a_1, \dots, a_\mu : R_1, \dots, R_\nu \rangle \\ \phi &\rightarrow F \\ a_i &\mapsto A_i \end{aligned}$$

where μ and ν are small. Here the A_i comprise a set of matrices generating F and the R_i comprise a complete set of relators. Each R_j is a word in the letters $a_i^{\pm 1}$. For each

word W in the letters $a_i^{\pm 1}$ we will use \bar{W} to denote the element of F which it represents. Thus $\bar{a}_i = A_i$.

It is clear that any crossed homomorphism is determined by its values on the generators, so Cr may be isomorphically embedded in $(Q/Z)^{n\mu}$ by $t \mapsto [t(A_1), t(A_2), \dots, t(A_\mu)]$. Let Cr', Pr' , denote the images of Cr, Pr , under this embedding.

Now we will use the relators to characterize Cr' and Pr' as subgroups of $(Q/Z)^{n\mu}$. We begin by defining a function

$$\begin{aligned} ((W, j) : W \text{ a word, } 1 \leq j \leq \mu) &\rightarrow \{n \times n \text{ integer matrices}\} \\ (W, j) &\longmapsto W^{(j)} \end{aligned}$$

with the property: if W is any word, and t is any crossed homomorphism, then $t(\bar{W}) = \sum_{j=1}^{\mu} W^{(j)} t(A_j)$.

One can verify that this property is enjoyed by the function:

$$\text{if } W = a_{i_1}^{\epsilon_1} \dots a_{i_p}^{\epsilon_p}, \quad \epsilon_j = \pm 1, p \geq 1,$$

$$\text{then } W^{(j)} = \sum_{\alpha=1}^p \delta_{i_\alpha j} \epsilon_\alpha \begin{matrix} \alpha - 1 & \epsilon_\beta \\ \Pi & A_{i_\beta} \\ \beta = 0 & A_j \end{matrix} (1/2)(\epsilon_\alpha - 1)$$

$$\text{and } (\text{empty})^{(j)} = 0,$$

wherein it is understood that A_{i_0} denotes the identity, and the δ is Kronecker's.

Another useful property of this function is:

$$\forall W, V, j, \quad (WV)^{(j)} = W^{(j)} + \bar{WV}^{(j)}$$

Now we form the $n\mu \times n\nu$ matrix

$$R = \begin{pmatrix} R_1^{(1)} & \dots & R_1^{(\mu)} \\ \cdot & & \cdot \\ \cdot & & \cdot \\ R_v^{(1)} & \dots & R_v^{(\mu)} \end{pmatrix}$$

=

R defines a group homomorphism $R : Q^{n\mu} \rightarrow Q^{n\nu}$ which maps the integer lattice $Z^{n\mu}$ into $Z^{n\nu}$, thus defining a homomorphism $R' : (Q/Z)^{n\mu} \rightarrow (Q/Z)^{n\nu}$ with the commuting diagram

$$\begin{array}{ccc} Q^{n\mu} & \xrightarrow{R} & Q^{n\nu} \\ \pi \downarrow & & \downarrow \pi \\ (Q/Z)^{n\mu} & \xrightarrow{R'} & (Q/Z)^{n\nu} \end{array}$$

Now, as Zassenhaus shows,

$$Cr' = \ker R'$$

$$Pr' = \pi \ker R.$$

Apply FTAG: $LRK = \text{diag} [e_1, e_2, \dots, e_\rho, 0, \dots, 0]$

Now, $Cr^* = K^{-1}Cr'$ is isomorphic to Cr' and

$$Cr^* = \ker R'K' = \ker L'R'K' \cong (Q/Z)^{n\mu-\rho} \oplus \left(\bigoplus_{i=1}^{\rho} Z_{e_i} \right)$$

$$\cup \quad \cup \quad \cup \quad \parallel$$

$$Pr^* = \pi \ker R'K' = \pi \ker LRK \cong (Q/Z)^{n\mu-\rho}$$

$$\text{Therefore, } H^1(F, (Q/Z)^n) \cong \bigoplus_{i=1}^{\rho} Z_{e_i}$$

may be represented by the elements

$$[f_1/e_1, f_2/e_2, \dots, f_\rho/e_\rho, 0, \dots, 0] \in Cr^*$$

where $0 \leq f_i < e_i$, for all i , $0 \leq i < \rho$.

Finally, each of these $n\mu$ -component column vectors is multiplied on the left by K and recognized as a crossed homomorphism

$$t: \{A_1, \dots, A_\mu\} \rightarrow (Q/Z)^n$$

which lifts to a function

$$s: \{A_1, \dots, A_\mu\} \rightarrow \mathbb{Q}^n$$

which in turn defines a crystallographic group G when the $s(A_i)$ are interpreted as the translation parts of a set of generators of G over \mathbb{Z}^n . Clearly, different lifts of a crossed homomorphism give the same (not merely equivalent) group G .

We will call the lift of a crossed homomorphism a Frobenius map, adapting nomenclature used by Burckhardt. Thus, a Frobenius map is any function $s: F \rightarrow \mathbb{Q}^n$ which satisfies the so-called Frobenius congruences,

$$\forall A, B \in F, \quad s(AB) \equiv s(A) + As(B) \pmod{\mathbb{Z}^n}.$$

We have called the computer code devised to execute this algorithm HILB18, commemorating Hilbert's 18th problem. (*)

The original presentation for F may be used together with the Frobenius map to obtain a presentation for G as follows:

Take for generators z_1, \dots, z_n , corresponding to the basis of \mathbb{Z}^n ,

and a_1, \dots, a_μ , corresponding to the generators of F .

Take for relators $z_i z_j z_i^{-1} z_j^{-1}$, $1 \leq i < j \leq n$,

and $a_i z_j a_i^{-1} z_1^{-\alpha(ij1)} \dots z_n^{-\alpha(ijn)}$, $1 \leq i \leq \mu$, $1 \leq j \leq n$,

and $R_j z_i^{-\beta(j1)} \dots z_n^{-\beta(jn)}$, $1 \leq j \leq \nu$,

(*) A similar program has been written by Janssen and Fast. [15].

where c_{ijk} is the (jk) component of A_i , R_j is a relator in the presentation for F , and $\beta(jk)$ is the k th component of the translation part of the element of G obtained by substituting $(A_i, s(A_i))$ for a_i in R_j .

We have not given these presentations in our tables of results because we have generally found the matrix representations more useful and suggestive, and the former are easily computed from the latter. We mention these presentations here only because we do use them in our computation of the first homology groups in Chapter 4.

Chapter 3

The Torsion-free Extensions

Our concrete treatment of crystallographic groups as subgroups of $GL(n, Z) \cdot Q^n$ allows a direct arithmetical solution to step III.

Definition: Let G be a crystallographic extension of F . Then we say that $A \in F$ is critical for G if G has a non-trivial element of finite order with linear part A .

If A is a matrix of finite order k , then we denote $\sum_{i=0}^{k-1} A^i$ by $\text{tr}A$. Clearly, $(I-A)\text{tr}A = 0$.

Proposition 1: If $(A, s) \in GL(n, Z) \cdot Q^n$, then (A, s) has finite order if and only if $\text{tr}As = 0$, and in this case the order of (A, s) is the order of A . Proof: Immediate from the multiplication rule.

Proposition 2: Let G be a crystallographic extension of F with Frobenius map s . Let $A \in F$ have order $k > 1$. Then A is critical for G if and only if the equation

$$(1) \quad \text{tr}A s(A) = \text{tr}A z$$

is satisfied by some $z \in Z^n$. Proof: Immediate from Proposition 1 and the fact that the elements of G with linear part A are just the $s(A) \cdot z$, $z \in Z^n$.

Proposition 2 gives the necessity of Condition A for the existence of torsion-free extensions of F , because the image of $\text{tr}A$ is contained in the fixed space of A and if this is zero then (1) is satisfied trivially independent of s . Now we can discover whether G is torsion-free by testing each $A \in F$ for

solvability of equation (1) by integers. This test is easily made mechanically as seen below. Note that the Frobenius congruences imply that $\text{tr} A s(A) \in \mathbb{Z}^n$.

The computing may be reduced considerably by noting that it suffices to apply the test to a certain proper subset of F which we proceed to define.

Proposition 3: Let C be conjugate to B in F , $C = ABA^{-1}$.

Then C is critical for G if and only if B is critical for G .

Proof: Let s be a Frobenius map for G . Then, by the Frobenius congruence, $s(C) = s(A) + As(B) - ABA^{-1}s(A) + z_1$

$$= (I-C)s(A) + As(B) + z_1, \text{ for some } z_1 \in \mathbb{Z}^n.$$

Now suppose B is critical. Then $\exists z_2 \in \mathbb{Z}^n$ with $\text{tr} Bs(B) = \text{tr} Bz_2$.

Then, since $\text{tr} C = A(\text{tr} B)A^{-1}$, we have

$$\begin{aligned} \text{tr} Cs(C) &= \text{tr} C((I-C)s(A) + As(B) + z_1) \\ &= \text{tr} C(As(B) + z_1) \\ &= A \text{tr} Bs(B) + \text{tr} Cz_1 \\ &= A \text{tr} Bz_2 + \text{tr} Cz_1 \\ &= \text{tr} C(Az_2 + z_1), \end{aligned}$$

showing that equation (1) is satisfied for C by $Az_2 + z_1 \in \mathbb{Z}^n$.

Proposition 4: If A , of order k , is critical for G , then so is any of its powers A^p with $1 \leq p \leq k - 1$. Proof: Trivial.

Proposition 5: Let $A \in F$ have order k and let s be a Frobenius map for G . If A^p is critical for G then there exists $z \in \mathbb{Z}^n$ satisfying

$$(2) \quad \text{ptr} A s(A) = \text{tr} A z.$$

Proof: By Proposition 4, we may assume, without loss of generality, that p divides k . Then $q = k/p$ is the order of A^p and

$$\text{tr}A^p = \sum_{i=0}^{q-1} A^{ip}.$$

$$\text{By Frobenius, } s(A^p) = \sum_{j=0}^{p-1} A^j s(A) + z_1, \quad z_1 \in \mathbb{Z}^n$$

$$\begin{aligned} \text{Whence } \text{tr}A^p s(A^p) &= \sum_{i=0}^{q-1} \sum_{j=0}^{p-1} A^{ip+j} s(A) + \text{tr}A^p z_1 \\ &= \text{tr}As(A) + \text{tr}A^p z_1. \end{aligned}$$

If A^p is critical, then $\text{tr}A^p s(A^p) = \text{tr}A^p z_2$, $z_2 \in \mathbb{Z}^n$,

$$\text{and } \text{tr}As(A) = \text{tr}A^p z_3, \quad z_3 = z_1 - z_2.$$

But the left side of this equation is fixed by A . Thus

$$p \text{tr}As(A) = \sum_{j=0}^{p-1} A^j \text{tr}As(A) = \sum_{j=0}^{p-1} A^j \text{tr}A^p z_3$$

$$= \text{tr}Az_3,$$

which is equation (2).

The converse to Proposition 5, even under the added hypothesis $p|k$, is not true. The simplest counterexample occurs in the torsion-free crystallographic extension of C_4 generated over \mathbb{Z}^4 by

$$(A, s) = \left(\begin{pmatrix} 0 & -1 & & \\ 1 & 0 & & \\ & & -1 & 1 \\ & & 0 & 1 \end{pmatrix}, \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1/2 \end{pmatrix} \right).$$

Here

$$\text{tr}As = (1/2) \text{tr}A \begin{pmatrix} 0 \\ 0 \\ 0 \\ 1 \end{pmatrix}$$

but A^2 is not critical.

Definition: Let ϕ be a finite group. Then we call a subset $T \subseteq \phi$ a torsion test set if every element of ϕ is conjugate to a power of some element of T .

Torsion test sets for all the abstract groups occurring in our initial sample of 4-dimensional crystal classes are given in Table B. Now let G be a particular crystallographic extension of F given by a Frobenius map s , and let T be a torsion test set for F . According to the foregoing propositions, if, for every $A \in T$ and every divisor p of the order of A , equation (2) has no solution by $z \in Z^n$, then G is torsion-free. If for some $A \in T$ equation (1) has a solution, then G is not torsion-free. If neither case obtains, and so for some $A \in T$ equation (2) has a solution with $1 < p|k$, then we must test A^p .

Deciding the solvability of (2) is another application of FTAG. Namely, if $L(\text{tr}A)K = \text{diag}(f_1, \dots, f_r, 0, \dots, 0)$ and $L\text{tr}A \cdot s(A) = [u_1, \dots, u_n]$, then the smallest positive integer p for which (2) has integral solution z is the least common multiple of the numbers $f_i / (f_i, u_i)$, $1 \leq i \leq r$.

Our computer code for Step III is called FINDFLAT.

Chapter 4

The First Homology Group and the Isomorphism Problem

We have implied in Chapter 0 that all of the affine structure of a compact euclidean space form M is contained in its fundamental (i.e., deck) group G . One quotient group of G , the holonomy G/\mathbb{Z}^n , has been instrumental for the construction of G . We now consider another quotient of G , its commutator quotient group $G/[G,G]$. This, of course, is isomorphic to the first integral homology group of M , $H_1(M, \mathbb{Z})$. In the present work it is an important tool in the completion of Step IV, the isomorphism problem. For $H_1(M, \mathbb{Z})$ is a finitely generated abelian group, and the isomorphism of two given finitely generated abelian groups is trivially decidable by FTAG; while to the best of the author's knowledge the isomorphism problem, in the sense of presentation theory, is unsolved for crystallographic groups.

Magnus, Karrass, and Solitar [20] give a mechanical procedure by which the commutator quotient group of any finitely presented group G may be determined from the presentation. We construct the so-called "titled exponent sum matrix" E . Its rows are labelled by the generators and its columns are labelled by the relators. The (ij) element is the sum of the exponents to which generator i appears in relator j . Apply FTAG: $LEK = \text{diag}(e_1, \dots, e_r, 0, \dots, 0)$. Now the e_i are the

torsion coefficients of $G/[G,G]$ and $\lambda-r$ is its Betti number, where λ is the number of generators in the presentation.

That is, $G/[G,G] \approx Z_{e_1} \oplus \dots \oplus Z_{e_r} \oplus Z^{\lambda-r}$

Our computer program ABELR carries this out for each of the groups determined in Step II according to the presentation given at the end of Chapter 2. The first homology groups of the 4-dimensional compact euclidean space forms are listed in Table D. We note in passing that there are at least two instances ([1], [5]) in the literature on compact euclidean space forms where the spaces with vanishing first Betti number present a special problem to the structure theory. Now $b_1 = 0$ for one of the ten 3-dimensional spaces, and Table D shows that in dimension 4 it happens in only 4 cases out of 73.

We turn now to the isomorphism problem itself. Let G_1 and G_2 be two crystallographic extensions of the same $F \subset GL(n, Z)$ with Frobenius maps s_1 and s_2 . Let k be the order of F and e_ρ the exponent of $H^2(F, Z^n)$, so that the $t_i = e_\rho s_i$ are functions $F \rightarrow Z^n$. By Bieberbach's theorem, G_1 and G_2 are isomorphic if and only if there is a matrix $X \in GL(n, Z)$ which normalizes F and for which the congruences

$$(1) \quad \forall A \in F \quad s_2(A) \equiv Xs_1(X^{-1}AX) + (A-I)u \pmod{Z^n}$$

have a solution $u \in Q^n$. Here the element $(X, s) \in GL(n, Z) \cdot Q^n$ conjugates G_1 to G_2 . Moreover, one can see from the proof in [26] (involving an "averaging over F ") that if a solution u exists then one may always be found in $(1/ke_\rho)Z^n$. Zassenhaus describes an efficient technique for finding all the isomorphisms

obtaining among the elements of $H^2(F, Z^n)$ if one has a finite set of generators of the normalizer of F in $GL(n, Z)$, and the author [17] has indicated an approach to the algorithmic construction of such a generating set. However, in the present computation we have taken a more direct, if less elegant, tack, because it proves entirely feasible for our problem and allows the use of several of the subroutines developed for the preceding steps.

The method is in the same spirit as our solution of the arithmetical equivalence problem of Step I, but rather more complicated. Once again, it suffices to consider the congruences (1) for all A in a set of generators $\{A_1, \dots, A_\mu\}$ of F . We introduce new indeterminates $z_i \in Z^n$, $i = 1, \dots, \mu$, converting the congruences into equations, which we then multiply by ke_ρ to obtain equations with integer coefficients. Thus, G_1 and G_2 are isomorphic if and only if, for some automorphism σ of F the equations

$$\begin{aligned} (a) \quad & A_i X - X \sigma(A_i) = 0 \\ (2)(b) \quad & kXt_1(\sigma(A_i)) + (A_i - I)u + ke_\rho z_i = kt_2(A_i) \quad i=1, \dots, \mu \\ (c) \quad & \det X = \pm 1 \end{aligned}$$

have a solution by an $n \times n$ unimodular integer matrix X and vectors $u, z_1, \dots, z_\mu \in Z^n$. For each σ (2ab) presents a system of $\mu(n^2 + 1)$ inhomogeneous linear equations in $n^2 + \mu n + n$ unknowns, and the left side of (2c) is an n -th degree homogeneous polynomial in the n^2 coefficients of X . Let $(2b)^\circ$ denote the homogeneous form of equation (2b), i.e. replace $kt_2(A_i)$ by 0. The comments of Chapter I restricting the automorphisms σ that we must consider apply here as well.

For fixed σ , G_1 , G_2 , all the matrices X belonging to solutions $\{X, u, z_1, \dots, z_u\}$ of (2ab) are represented by a matrix Ξ whose components are inhomogeneous linear forms in some indeterminates ξ_1, \dots, ξ_r , and the question of the isomorphism of G_1 and G_2 is reduced to the question of the existence of integer solutions to one of the polynomial equations $\det \Xi = \pm 1$ for some admissible σ . Once again, the linear algebra and the computation of the polynomial proceed automatically and the latter is printed out only if it satisfies some obvious necessary conditions. Deciding the existence of integer roots still proceeds smoothly.

We have not applied this method to find all the isomorphism classes of crystallographic extensions of F , but only to decide the isomorphisms for the torsion-free extensions, and these were first divided into classes of common commutator quotient groups. Further, in the program we do not solve (2ab) in one step as implied above, but in three steps, solving first (2a), then (2b) $^\circ$, then finally (2b). This breakdown is made because (2a) depends only on σ and therefore must be solved only once for each σ admissible for F ; and (2b) $^\circ$ depends only on σ and the group G_1 , so that if there are m groups in one of our classes, then (2b) $^\circ$ must be solved at most $m-1$ times for each σ , while (2b) must be solved $m(m-1)/2$ times; and further, the full matrix of coefficients of (2ab) was sometimes too big to work with in the machine memory available to us.

In order to justify the code listed in Table E, we will show in a little greater detail how this procedure has been

reduced to the mechanical manipulation of integer matrices.

We construct the $\mu n^2 \times n^2$ matrix of coefficients M_σ of (2a), and regard it as a map $M_\sigma: Z^{n^2} \rightarrow Z^{\mu n^2}$ whose kernel is the solution set of (2a) and hence contains all X represented by Ξ . Apply FTAG: $L_\sigma M_\sigma K_\sigma = \text{diag}(e_1, \dots, e_{r_\sigma}, 0, \dots, 0)$. Let $d_\sigma = n^2 - r_\sigma$ and let K'_σ be the rightmost d_σ columns of K_σ ; these columns form an integral basis for $\ker M_\sigma$. That is, we have expressed the general matrix solution to (2a) as a matrix of homogeneous linear forms in some variables w_1, \dots, w_{d_σ} .

Substitution of these forms for the components of X in (2b) makes the latter into a system of μn inhomogeneous linear equations in $d_\sigma + \mu n + n$ unknowns. Let M_b be the matrix of coefficients of this system, and regard it as a map $M_b: Z^{d_\sigma} \oplus Z^{\mu n + n} \rightarrow Z^{\mu n}$. We have decomposed the domain in order to emphasize that the leftmost d_σ columns of M_b contain the coefficients of the w_i in (2b) and that we are ultimately not interested in the values, but rather only the existence of, the u, z_i , in the solutions. Now $\ker M_b$ is the solution set of (2b) $^\circ$. Apply FTAG: $L_b M_b K_b = \text{diag}(f_1, \dots, f_{r_b}, 0, \dots, 0)$. Let $d_b = d_\sigma + \mu n + n - r_b$ and let K''_b be the upper left $d_\sigma \times r_b$ corner and K'_b the upper right $d_\sigma \times d_b$ corner of K_b . The rightmost d_b columns of K_b form an integral basis for $\ker M_b$ and the columns of K'_b generate $\text{proj}_{d_\sigma}(\ker M_b)$ over Z . But we need a basis for the projection, so apply FTAG. $L_c K'_b K_c = \text{diag}(g_1, \dots, g_{r_c}, 0, \dots, 0)$. Let L'_c be the leftmost r_c columns of $K'_b K_c$. Then the columns of

$K'_\sigma L'_c$ form an integral basis for the set of all X in Z which satisfy (2a) and for which there are solutions of (2b) $^\circ$. All these, and only these, are represented by the matrix Ξ° of homogeneous linear forms given by:

$$\Xi_{ij}^\circ = \sum_{k=1}^{r_c} (K'_\sigma L'_c)_{ck} \xi_k,$$

where $c = i + (j-1)n$ is the index of the column of M_σ containing the coefficients of the (ij) component of X in (2a).

The solution set of (2b) is a coset of $\ker M_b$, so we need to find only one solution. Write the right side of (2b) as a vector $\bar{v} = k[t_2(A_1), \dots, t_2(A_u)] = [v_1, \dots, v_{un}]$. Now one can check that there is no solution and so no isomorphism $G_1 \approx G_2$ for σ , unless

$$\begin{aligned} f_i \mid (L_b \bar{v})_i & \quad 1 \leq i \leq r_b \\ (L_b \bar{v})_i = 0 & \quad r_b < i \leq un, \end{aligned}$$

in which case a solution is given by

$$\bar{u} = [(L_b \bar{t})_1 / f_1, \dots, (L_b \bar{t})_{r_b} / f_{r_b}, 0, \dots, 0].$$

Finally,

$$\Xi_{ij} = \Xi_{ij}^\circ + (K'_\sigma K''_b \bar{u})_a, \quad c = i + (j-1)n,$$

gives an $n \times n$ matrix of inhomogeneous linear forms in the variables ξ_k , such that the existence of integer solutions of one of the polynomial equations $\det \Xi = \pm 1$ is equivalent to the existence of an isomorphism $G_1 \approx G_2$ inducing σ on F .

The necessary condition:

gcd {coefficients of $\det \bar{E}$ } = 1

and gcd {coefficients of non-constant terms}

divides one of (coefficient of constant term) ± 1

is tested before printing out the polynomial.

The foregoing constitutes a fairly faithful description of the flow in the core parts of the programs ISG/ISH.

Chapter 5

The Computer Programs

The mechanical parts of this computation were accomplished on the Control Data 3150 computer at Humboldt State College. This machine has a core storage of 2^{14} 24-bit words, of which about 8000 are available to the programmer over the operating system and system I/O routines. Available memory was the principal constraint posed by the machine. Nevertheless, the running versions of the codes have sufficient storage allocated to carry the computations into dimension 5 for the groups which can be generated by 3 elements. This was achieved by using a high degree of segmentation of functions; i.e., we have many small main programs, each doing a small part of the work on one pass over all the data, with intermediate results stored on magnetic tape. On the other hand, these main programs share to a high degree their workhorse subroutines.

The constraint posed by the machine's speed (3.5 microsecond storage access time) was negligible. The 16K operating system does not include run timing, so we have no data on the speed of the programs. Suffice it to say that run time was always small when compared both with turnaround time and machine idle time.

The effect of the hardware constraint of size was intensified by the existence of a countermanding psychological constraint, viz. the desire to keep the programs simple and therefore finitely debuggable and ultimately believable. The coding for the most

part is in Fortran [28], with the few exceptional subroutines coded in the Control Data assembly language Compass.^[27] In Table E we give the source language listings of the actual versions used to obtain our main result.

It will be noted that there is a fair amount of redundancy in our computations. For example, in some places we will carry out the arithmetic to test whether A is equivalent to C, ignoring the fact that we have already established, say, both that A is equivalent to B and that B and C are not equivalent. These redundancies originated in our deeming too dear the cost in program complexity of their removal. However, they proved to be a definite advantage when noted inconsistencies among redundant results revealed the presence of subtle and long undetected bugs. And the ultimately achieved complete consistency among all the redundancies adds significantly to our confidence in the accuracy of our results.

Almost all of the codes were thoroughly tested on the analogous and previously known 3-dimensional cases before application to our main problem. But we know, by the bugs undetected in this phase and later discovered through inconsistent 4-dimensional results, that the 3-dimensional test did not explore all of the logical loops and branches.

We conclude this paper with a few remarks on each of the listed programs to aid the reader who wishes to study them. We begin with the Compass codes.

Subroutine DIAR.

This is our implementation of the Fundamental Theorem of

Generated Abelian Groups, FTAG. The algorithm is a slight modification of the constructive proof in [13]. We have coded it in assembly language because it is our main workhorse and it is long, whereas our Fortran compiler is particularly space-wasteful.

The main entry-point calling sequence is CALL DIAKL
(NL,ML,NK,MK,NM,MM, KK, NR, NC).

The inputs are: MM, an arbitrary NR x NC integer matrix;
NM, NL, NK, the calling program row-dimensions of the matrices
MM, ML, MK (resp.).

The outputs are: MM, the diagonal form of the input MM as described in Chapter 0; ML and MK, the left and right transforming unimodular matrices; KK, the number of non-zero diagonal elements.

The transforming matrices are computed by beginning with identity matrices and performing at each "pivot step" elementary row and column operations corresponding to those applied to MM.

There are two other entry-points, DIAK, DIALI, each with an appropriately truncated calling sequence. The former omits the computation of the left transforming matrix and the latter omits the computations of the right transforming matrix while computing the inverse of the left transforming matrix.

DIAR is called actually by only one of the main programs, ISG, which is the only one which requires all of the functions in core at one time. The identifiers DIAGK, DIAGL, DIAG, etc., appearing in the remaining main programs, are entry-points to other versions which differ from DIAR only in omitting one or more of its functions with respect to the transforming matrices.

We have omitted these shorter versions from Table E.

Subroutine MMPYR.

This is our integer matrix multiplier. The main entry-point calling sequence `CALL MMPYI(KA,KB,KC,N1,N2,N3,O,MA,MB,MC)` multiplies the $N1 \times N2$ matrix MA by the $N2 \times N3$ matrix MB, returning the product in MC. KA, KB and KC are the calling-program row dimensions of these matrices. The alternate entry points enable one to set in the subroutine some of the values or addresses of the input parameters, and to enter without passing some input values or addresses if unchanged since the previous call.

Function IGCDR.

`IGCD(N,M)` has the value of the greatest common divisor of the input integers N,M; Euclid's algorithm coded in Compass in order to exploit a certain feature of the hardware; viz., on a fixed-point divide, the quotient is left in one register (A), while the remainder is left in another (Q).

Subroutine POLYR.

This is our co-routine for manipulating integer polynomials in several variables. It is used with Program VC in Step I and with Program ISH in Step IV.

A polynomial is represented by a linked list. Each term of the polynomial is represented by a cell in the list. A cell consists of three consecutive words of core storage; but the cells representing consecutive terms of the polynomial are not necessarily consecutive in core. The first word of a cell contains the absolute core address of the cell representing the succeeding term, zero in the cell representing the last term. The calling Fortran program knows the polynomial by an integer identifier

naming a core location which contains the absolute core address of the cell representing the first term.

The second word in each cell gives the variable part of the term. Polynomials in up to 23 variables are allowed; each variable is coded by an integer between 1 and 23 and zero codes the constant term. The product of variables comprising the variable part of the term is coded by writing the variables in a string with the higher variables to the right and repetitions denoting powers, and then taking the value of this string read as a base-24 numeral. Thus, with our 24-bit word, POLYR is limited to polynomials of degree ≤ 5 .

The third word in each cell contains the integer coefficient of the term.

A single block of storage, STACK + 1 to STACK + 999, is allocated for all the polynomial lists. At any time, all cells not currently in use by some polynomial are linked together through their first words into one list of free cells which operates as a pushdown stack. Location STACK contains the absolute core address of the first free cell.

We now describe the functions by calling sequence.

CALL CLSTK.

This erases all existing polynomials by linking all the cells of allocated storage into the stack of free cells.

CALL CRPHDI(N,MV,MC,KP)

This creates a list representing the 1st degree polynomial with N terms whose variables (integers between 1 and 23) are given in the array MV and whose coefficients are given in the

array MC. On return, KP contains the absolute core address of the first cell of the list.

CALL PDCOD(KP,MV,MC,NT,NTM)

This supplies to the arrays MV and MC of the calling program the variable-part codes and coefficients of the polynomial KP. The output parameter NT gives the number of terms, while the input parameter NTM specifies the maximum number of terms which the calling program can accept.

CALL FREPOL(KP)

erases polynomial KP by attaching its tail to the head of the stack.

CALL PML1(KP,KQ,KR, ± 1).

According to the sign of the fourth parameter, this replaces polynomial KR by $KR \pm KP \cdot KQ$, where KP is assumed to be of degree ≤ 1 .

CALL UNHMG(KP,MC)

This assumes that KP is a polynomial without constant term, and adds to it, as leading term, the constant MC

CALL HOMOG(KP)

removes from KP its first term provided that the latter is constant.

Now we go on to the Fortran codes, beginning with the more general subroutines.

Function IDET4(M)

has as its value the (absolute core location of) 4th degree polynomial defined as the determinant of the 4 x 4 input matrix M whose components are polynomials of degree ≤ 1 .

Subroutine NITTY (NA,NB,MU,MA,MB,MDEL).

The first five parameters are input. They specify a set of MU $NA \times NA$ matrices stored in the array MA and a set of MU $NB \times NB$ matrices contained in the array MB. Write the matrix equations

$$(1) \quad MA_i X - X MB_i = 0 \quad i = 1, \dots, MU$$

Where X is an indeterminate $NA \times NB$ matrix. Then the output MDEL is the $(NANBMU) \times (NANB)$ matrix of coefficients of equations (1) considered as a system of $NANBMU$ equations in the $NANB$ unknown components X_{jk} of X. The coefficients of the (jk) component of the ith equation of (1) are contained in row $j + (k-1)NA + (i-1)NANB$ of MDEL; the coefficients of the variable X_{jk} are contained in column $j + (k-1)NA$. NITTY is called by Programs VB, VC and ISG. Our Fortran system requires agreement between calling program and subroutine in the row dimension of matrix parameters, so several versions of NITTY, differing only in their DIMENSION statements, are used. Only one version is listed in Table E. Similar remarks are applicable to all of our Fortran matrix-manipulation subroutines and will not be repeated.

Subroutine MATP (ND, NP, MA, MB)

returns to MB the NPth power of the $ND \times ND$ matrix MA. The alternate entry-point MATT returns to MB instead the sum of the powers 0 through NP of MA. MATT/MATP is called by Subroutine WORDC and by Program FINDFLAT.

Subroutine WORDC (ND, LI, LE, MG, MGA)

The inputs are: ND, the dimension; (LI, LE), a word as explained under Program ACCHK; MG, an array containing a set of

NDXND matrices. The output is MGA, the matrix obtained by substituting the matrices of MG for the corresponding letters of the word. Called by VC.

Subroutine WORDV(LI,LE,MX,MI).

The input is a word W specified by the arrays (LI,LE) as explained under Program ACCHK. The output array MX contains the matrix values $w^{(j)}$, $1 \leq j \leq MU$, of the function defined in Chapter 2. The generating matrices are obtained through COMMON storage. The output MI is the matrix value of w . WORDV is called by Programs ACCHK and ISG.

We proceed now to the main programs. Each of these produces copious printed output displaying the input data as well as the computed results, and we will generally not mention this printout in the I/O comments in each case.

Program VB.

This program carries out the first part of Step I, the computation of the cohomology group of binding functions. The input is a deck of cards consisting of a subdeck for each geometrical crystal class. The first card of a subdeck carries a few parameters identifying and describing the class. Then follows a set of NAUT automorphisms of the class to be used by Program VC in deciding arithmetical equivalence as described in Chapter 1. Each automorphism is specified by giving the images of the generators as words in the generators. See under Program ACCHK for the coding of words. Then follows 2·NDEC·MU cards carrying matrix representatives of the generators, thus defining NDEC (NA+NB)-decomposed representations of the class. The components

of these are stored in the arrays MA and MB. The cohomology group of binding functions is computed as described in Chapter 1. In addition to the printout, the output consists of

(i) punched cards carrying the representations of degree $ND = NA + NB$ defined by all the elements of all the cohomology groups computed for the geometrical class. These are in the proper format for input to Program ACCHK, but of course contain perhaps several representatives of each arithmetical crystal class. It is at this point that the arithmetical classes acquire the "AC" designations of Table C.

(ii) A scratch tape carrying all the input data and computed results for immediate use by the next program.

Program VC

This reads through the scratch tape just written by VB. For each geometrical crystal class and each pair of representatives of arithmetical classes generated by VB, and each of the given automorphisms, VC computes the polynomial $\det \Xi$ according to the scheme of Chapter 1, printing it out if it satisfies the necessary condition mentioned. After human verification of the existence of integer roots, the redundant representatives are removed from the deck punched by VB.

Program ACCHK.

This code performs some editorial pre-processing on the data to be used by the programs of Steps II, III, and IV. The input is a deck of cards comprised of a number of subdecks each pertaining to an isomorphism class of point groups. The first card of a subdeck carries identifying information and some

parameters of the group, e.g. number of generators, number of relators, etc. Then comes a set of cards giving the relators of a presentation for the group. Each relator is a word in the generators and is coded as a sequence of up to 4 pairs (LRI,LRE) of positive integers; each pair designates a letter in the word, LRI naming the generator and LRE giving the exponent. The relators are followed by cards carrying the elements of a torsion test set, specified also as words in the generators, and also the orders of these elements. The remainder of the isomorphism class subdeck consists of the cards punched out by Program VB giving arithmetical crystal classes of the same isomorphism class.

In each arithmetical class, ACCHK begins by verifying that the input generating matrices indeed satisfy the given relators. During this test the matrix R of Chapter 2, designated MR in the code, is produced. Incidentally, this matrix has the property that, if multiplied on the right by the (nu) -component column vector formed from the values on the generators of a Frobenius map for an extension of the class, then the components of the product are just the $\beta(jk)$ of the presentation for the extension as defined in Chapter 2.

Next, ACCHK computes, from the given generating matrices and torsion test words, the corresponding matrix representatives of the torsion test elements (array MT), and verifies that these matrices have the given orders.

Finally, the program constructs the matrix MRT with the property that multiplying it by the vector containing the values on the generators of a Frobenius map yields the values on

the torsion test elements of a Frobenius map for the same extension.

All of the input and computed result are written on a scratch tape for use by the next program.

Program HILB18.

The computations of $H^2(F, Z^n)$ now consists merely of diagonalizing the matrix MR found on the tape written by ACCHK.

The output tape produced by HILB18 is saved as the main source of data for Steps III and IV. This tape contains several Fortran logical records for each arithmetical crystal class. The first record for each class contains almost all the information on the class so far accumulated; including the parameters of $H^2(F, Z^n)$. Then for each nonzero element of the latter there is a short record carrying certain parameters of the corresponding extension: identifiers; values of a Frobenius map on the generators and torsion test elements, and the $\beta(jk)$ of the presentation. The code provides the ability to update this tape.

Program FINDFLAT:

This determines, according to the scheme of Chapter 3, whether each crystallographic extension represented on the HILB18 output tape is torsion-free. In each arithmetical class and for each torsion test element A, the computation of $\text{tr}A$ and its diagonalization are not performed until needed, and then are saved should they be needed again. However, for the p-th power of a torsion test element the same computations are carried out anew each time they are needed.

Program ABELR.

computes the commutator quotient group of each extension on the HILB18 output tape, as described in Chapter 4.

Program ISG:

The input is the HILB18 output tape plus a deck of cards specifying sets of extensions to be tested for isomorphism and the admissible automorphisms for the arithmetical crystal classes involved.

Until this point in the computation, no attention was paid to the possible choices (mod Z^n) for the Frobenius maps. But the polynomials produced in Step IV were in some cases rendered more tractable by beginning with Frobenius maps with small components; so, ISG has an (internal) subroutine which chooses the Frobenius map with all of its values on the generators in the cube $0 \leq s_i < 1$, $i = 1, 2, 3, 4$.

The core part of this code was detailed in Chapter 4. In each instance of equations (2ab) the code reports in the printout whether solutions exist; when they do, then the matrix Ξ of inhomogeneous linear forms is saved on a scratch tape for

Program ISH

whose function is analogous to that of Program VC in Step I.

The short editing codes which produced Tables C and D from the data tapes are omitted from the listings in Table E.

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APPENDIX**Tables**

Table A

This table displays the 45 4-dimensional geometrical crystal classes which satisfy Condition A.

Column G: The symbol G_{Ck} , $0 \leq k \leq 44$, is an arbitrary name we have given the class for reference. Classes G_{C0} to G_{C10} are orientation preserving; G_{C11} to G_{C44} are not.

Column I: The abstract group. C_k cyclic of order k
 D_k dihedral of order $2k$
 T tetrahedral
 θ octahedral

Column A: The number of arithmetical classes into which the geometrical class splits.

Column H: The number of the latter which occur as holonomy classes of compact euclidean space forms.

Column Hurley: The very useful notation of [11, 12, 18] which gives the distribution of elements in the conjugate classes of $GL(4, Q)$. The latter which contain crystal classes are characterized by the order k and coefficients χ , σ , d , of the characteristic equation: $\lambda^4 - \chi\lambda^3 + \sigma\lambda^2 - \chi\lambda + d = 0$

Here is a short dictionary for those with non-zero fixed space:

Subtable Hurley

	I	E	F	K	N	N'	R	T	T'	Z	
k	1	2	4	3	6	6	4	2	2	6	order
χ	4	0	0	1	1	-1	2	2	-2	3	trace
σ	6	-2	0	0	0	0	2	0	0	4	second invariant
d	1	1	-1	1	-1	-1	1	-1	-1	1	determinant

Table A

<u>G</u>	<u>I</u>	<u>A</u>	<u>H</u>	<u>Hurley</u>
GC0	C_1	1	1	1I
GC1	C_2	3	2	1I 1E
GC2	$C_2 \times C_2$	13	6	1I 3E
GC3	C_4	2	2	1I 1E 2R
GC4	D_4	6	2	1I 5E 2R
GC5	C_3	2	2	1I 2K
GC6	C_6	1	1	1I 1E 2K 2Z
GC7	D_3	8	3	1I 3E 2K
GC8	D_6	2	1	1I 7E 2K 2Z
GC9	T	6	2	1I 3E 8K
GC10	\emptyset	5	0	1I 9E 8K 6R
GC11	C_2	2	2	1I 1T
GC12	C_2	2	1	1I 1T'
GC13	$C_2 \times C_2$	6	5	1I 1E 2T
GC14	$C_2 \times C_2$	7	2	1I 1E 1T 1T'
GC15	$C_2 \times C_2$	6	1	1I 1E 2T'
GC16	$C_2 \times C_2 \times C_2$	12	1	1I 3E 3T 1T'
GC17	$C_2 \times C_2 \times C_2$	12	1	1I 3E 1T 3T'
GC18	C_4	7	4	1I 1E 2F
GC19	$C_4 \times C_2$	6	2	1I 1E 2F 2R 1T 1T'
GC20	D_4	2	0	1I 1E 2R 4T

Table A (Page 2)

<u>G</u>	<u>I</u>	<u>A</u>	<u>H</u>	<u>Hurley</u>
GC21	D_4	2	0	1I 1E 2R 4T'
GC22	D_4	13	1	1I 3E 2F 2T
GC23	D_4	13	1	1I 3E 2F 2T'
GC24	$D_4 \times C_2$	6	0	1I 5E 2F 2R 5T 1T'
GC25	$D_4 \times C_2$	6	0	1I 5E 2F 2R 1T 5T'
GC26	D_3	3	0	1I 2K 3T
GC27	D_3	3	0	1I 2K 3T'
GC28	C_6	4	1	1I 2K 2N'1T
GC29	C_6	4	2	1I 2K 2N 1T'
GC30	$C_6 \times C_2$	2	1	1I 1E 2K 2N 2N'1T 1T'2Z
GC31	D_6	1	0	1I 1E 2K 6T 2Z
GC32	D_6	1	0	1I 1E 2K 6T'2Z
GC33	D_6	6	0	1I 3E 2K 2N'4T
GC34	D_6	6	0	1I 3E 2K 2N 4T'
GC35	D_6	6	0	1I 3E 2K 2N 3T 1T'
GC36	D_6	6	0	1I 3E 2K 2N 1T 3T'
GC37	$D_6 \times C_2$	2	0	1I 7E 2K 2N 2N'7T 1T'2Z
GC38	$D_6 \times C_2$	2	0	1I 7E 2K 2N 2N'1T 7T'2Z
GC39	$T \times C_2$	5	0	1I 3E 8K 8N'1T 3T'
GC40	$T \times C_2$	5	1	1I 3E 8K 8N 3T 1T'
GC41	\emptyset	6	0	1I 3E 6F 8K 6T'

Table A (Page 3)

<u>G</u>	<u>I</u>	<u>A</u>	<u>H</u>	<u>Hurley</u>
GC42	∅	6	0	1I 3E 6F 8K 6T
GC43	∅xC ₂	5	0	1I 9E 6F 8K 8N'6R 1T 9T'
GC44	∅xC ₂	5	0	1I 9E 6F 8K 8N 6R 9T 1T'

Table B

This table displays some information input to the programs which depends only on the isomorphism class of the point group; namely, presentations and torsion test sets.

The column headed NG gives the number of generators for the presentation. All the groups in question are generated by three or fewer elements, and these are designated A,B,C.

The column headed Relators gives a complete set of relators for the presentation.

The column headed Torsion Test gives a torsion test set as defined in Chapter 3 and as used in the computations.

Table B

<u>Group</u>	<u>NG</u>	<u>Relators</u>	<u>Torsion Test</u>
C_2	1	A^2	A
$C_2 \times C_2$	2	$A^2, B^2, (AB)^2$	A, B, AB
$C_2 \times C_2 \times C_2$	3	A^2, B^2, C^2 $(AB)^2, (AC)^2, (BC)^2$	A, B, C, AB, AC BC, ABC
C_4	1	A^4	A
D_4	2	$A^4, B^2, (AB)^2$	A, B, AB
$C_4 \times C_2$	2	A^4, B^2, A^3BAB	A, B, AB, A^2B
$D_4 \times C_2$	3	$A^4, B^2, C^2, (AB)^2$ $A^3CAC, (BC)^2$	A, B, C, AB AC, BC, A^2C
C_3	1	A^3	A
D_3	2	$A^3, B^2, (AB)^2$	A, B
C_6	1	A^6	A
D_6	2	$A^6, B^2, (AB)^2$	A, B, AB
$C_6 \times C_2$	2	A^6, B^2, A^5BAB	A, AB, A^2B
$D_6 \times C_2$	3	$A^6, B^2, C^2, (AB)^2$ $A^5CAC, (BC)^2$	A, B, AB, AC A^2C, BC, ABC
T	2	$A^3, B^3, (AB)^2$	A, B, AB
$T \times C_2$	3	$A^3, B^3, C^2, (AB)^2,$ A^2CAC, B^2CBC	AB, AC, ABC
\emptyset	2	$A^3, B^4, (AB)^2$	A, B, AB

Table C

This table displays the 49 4-dimensional arithmetical classes of holonomy.

The first, unheaded, column contains 3 entries for each class. The first, $AC\ k$, is an arbitrary designator used for reference on our data tapes. The second and third give the geometrical class and isomorphism class to which the arithmetical class belongs, in the same notation as Table A.

The 3 columns headed $GEN\ A,B,C$ give matrix representatives of a set of generators for the class; these correspond to the generators A,B,C of the presentations of Table B.

The last column headed $NFLAT$ gives the number of inequivalent compact euclidean space forms of the class.

TABLE C P. 1

	GEN A				GEN B				GEN C				NFLAT
AC 0	1	0	0	0									1
GC0	0	1	0	0									
C1	0	0	1	0									
	0	0	0	1									
AC 1	-1	0	0	0									1
GC1	0	-1	0	0									
C2	0	0	1	0									
	0	0	0	1									
AC 2	-1	0	0	1									1
GC1	0	-1	0	0									
C2	0	0	1	0									
	0	0	0	1									
AC 7	-1	0	0	0	1	0	0	0					4
GC2	0	-1	0	0	0	-1	0	0					
C2XC2	0	0	1	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC 8	-1	0	0	1	1	0	0	0					1
GC2	0	-1	0	0	0	-1	0	0					
C2XC2	0	0	1	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC 10	-1	0	0	1	1	0	0	0					1
GC2	0	-1	0	1	0	-1	0	1					
C2XC2	0	0	1	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC 14	-1	0	0	1	1	0	0	0					1
GC2	0	-1	0	1	0	-1	0	1					
C2XC2	0	0	1	0	0	0	-1	1					
	0	0	0	1	0	0	0	1					

TABLE C P. 2

	GEN A				GEN B				GEN C	NFLAT
AC 23	-1	0	0	0	0	1	0	0		1
GC2	0	-1	0	0	1	0	0	0		
C2XC2	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 24	-1	0	0	1	0	1	0	0		1
GC2	0	-1	0	1	1	0	0	0		
C2XC2	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 27	0	1	0	0						1
GC3	-1	0	0	0						
C4	0	0	1	0						
	0	0	0	1						
AC 28	0	1	0	0						1
GC3	-1	0	0	1						
C4	0	0	1	0						
	0	0	0	1						
AC 30	0	1	0	0	0	1	0	0		3
GC4	-1	0	0	0	1	0	0	0		
D4	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 31	0	1	0	0	0	1	0	0		1
GC4	-1	0	0	1	1	0	0	0		
D4	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 36	0	-1	0	0						1
GC5	1	-1	0	0						
C3	0	0	1	0						
	0	0	0	1						

TABLE C P. 3

	GEN A				GEN B				GEN C	NFLAT
AC 37	0	-1	0	0						1
GC5	1	-1	0	1						
C3	0	0	1	0						
	0	0	0	1						
AC 40	0	1	0	0						1
GC6	-1	1	0	0						
C6	0	0	1	0						
	0	0	0	1						
AC 41	0	-1	0	0	0	1	0	0		1
GC7	1	-1	0	0	1	0	0	0		
D3	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 45	0	-1	0	0	0	-1	0	0		1
GC7	1	-1	0	0	-1	0	0	0		
D3	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 47	0	-1	0	0	0	-1	0	0		1
GC7	1	-1	0	2	-1	0	0	0		
D3	0	0	1	0	0	0	-1	2		
	0	0	0	1	0	0	0	1		
AC 51	0	1	0	0	0	1	0	0		1
GC8	-1	1	0	0	1	0	0	0		
D6	0	0	1	0	0	0	-1	0		
	0	0	0	1	0	0	0	1		
AC 53	0	0	1	0	0	-1	0	0		1
GC9	1	0	0	0	0	0	-1	0		
T	0	1	0	0	1	0	0	0		
	0	0	0	1	0	0	0	1		

TABLE C P. 4

	GEN A				GEN B				GEN C				NFLAT
AC 54	0	0	1	0	0	-1	0	1					1
GC9	1	0	0	0	0	0	-1	1					
T	0	1	0	0	1	0	0	0					
	0	0	0	1	0	0	0	1					
AC 66	1	0	0	0									1
GC11	0	1	0	0									
C2	0	0	-1	0									
	0	0	0	1									
AC 67	1	0	0	0									1
GC11	0	1	0	0									
C2	0	0	-1	1									
	0	0	0	1									
AC 69	-1	0	0	0									1
GC12	0	-1	0	0									
C2	0	0	-1	0									
	0	0	0	1									
AC 77	-1	0	0	0	-1	0	0	0					6
GC13	0	-1	0	0	0	1	0	0					
C2XC2	0	0	1	0	0	0	1	0					
	0	0	0	1	0	0	0	1					
AC 78	-1	0	0	1	-1	0	0	1					3
GC13	0	-1	0	0	0	1	0	0					
C2XC2	0	0	1	0	0	0	1	0					
	0	0	0	1	0	0	0	1					
AC 80	-1	0	0	1	-1	0	0	1					1
GC13	0	-1	0	1	0	1	0	0					
C2XC2	0	0	1	0	0	0	1	0					
	0	0	0	1	0	0	0	1					

TABLE C P. 5

	GEN A				GEN B				GEN C				NFLAT
AC 81	0	1	0	0	1	1	1	0					1
GC13	1	0	0	0	0	0	-1	0					
C2XC2	-1	-1	-1	0	0	-1	0	0					
	0	0	0	1	0	0	0	1					
AC 84	-1	0	0	0	0	-1	0	0					1
GC13	0	-1	0	0	-1	0	0	0					
C2XC2	0	0	1	0	0	0	1	0					
	0	0	0	1	0	0	0	1					
AC 88	-1	0	0	0	-1	0	0	0					4
GC14	0	-1	0	0	0	-1	0	0					
C2XC2	0	0	1	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC 96	-1	0	0	0	-1	0	0	0					1
GC14	0	0	-1	0	0	-1	0	0					
C2XC2	0	-1	0	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC100	-1	0	0	0	-1	0	0	0					1
GC15	0	-1	0	0	0	-1	0	0					
C2XC2	0	0	1	0	0	0	-1	0					
	0	0	0	-1	0	0	0	1					
AC112	-1	0	0	0	1	0	0	0	-1	0	0	0	10
GC16	0	-1	0	0	0	-1	0	0	0	-1	0	0	
C2XC2XC2	0	0	1	0	0	0	-1	0	0	0	-1	0	
	0	0	0	1	0	0	0	1	0	0	0	1	
AC130	-1	0	0	0	1	0	0	0	-1	0	0	0	2
GC17	0	-1	0	0	0	-1	0	0	0	-1	0	0	
C2XC2XC2	0	0	1	0	0	0	-1	0	0	0	-1	0	
	0	0	0	-1	0	0	0	-1	0	0	0	1	

TABLE C P. 6

	GEN A				GEN B				GEN C				NFLAT
AC148	0	-1	0	0									1
GC18	1	0	0	0									
C4	0	0	-1	0									
	0	0	0	1									
AC150	0	-1	0	0									1
GC18	1	0	0	0									
C4	0	0	-1	1									
	0	0	0	1									
AC151	0	-1	0	0									1
GC18	1	0	0	1									
C4	0	0	-1	1									
	0	0	0	1									
AC152	-1	0	1	0									1
GC18	-1	0	0	0									
C4	-1	1	0	0									
	0	0	0	1									
AC156	0	1	0	0	-1	0	0	0					2
GC19	-1	0	0	0	0	-1	0	0					
C4XC2	0	0	1	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC160	1	0	-1	0	-1	0	0	0					1
GC19	1	0	0	0	0	-1	0	0					
C4XC2	1	-1	0	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC170	0	1	0	0	0	-1	0	0					1
GC22	-1	0	0	0	-1	0	0	0					
D4	0	0	1	1	0	0	1	0					
	0	0	0	-1	0	0	0	1					

TABLE C P. 7

	GEN A				GEN B				GEN C				NFLAT
AC171	0	1	0	0	0	-1	0	0					1
GC22	-1	0	0	1	-1	0	0	0					
D4	0	0	1	1	0	0	1	0					
	0	0	0	-1	0	0	0	1					
AC176	0	1	0	0	0	1	0	0					1
GC23	-1	0	0	0	1	0	0	0					
D4	0	0	1	0	0	0	-1	0					
	0	0	0	-1	0	0	0	1					
AC209	0	-1	0	0									1
GC28	1	-1	0	0									
C6	0	0	-1	0									
	0	0	0	1									
AC215	0	0	-1	0									1
GC29	-1	0	0	0									
C6	0	-1	0	0									
	0	0	0	1									
AC217	0	1	0	0									1
GC29	-1	1	0	0									
C6	0	0	-1	0									
	0	0	0	1									
AC219	0	1	0	0	-1	0	0	0					1
GC30	-1	1	0	0	0	-1	0	0					
C6XC2	0	0	1	0	0	0	-1	0					
	0	0	0	1	0	0	0	1					
AC257	0	0	1	0	0	-1	0	0	-1	0	0	0	1
GC40	1	0	0	0	0	0	-1	0	0	-1	0	0	
TXC2	0	1	0	0	1	0	0	0	0	0	-1	0	
	0	0	0	1	0	0	0	1	0	0	0	1	

Table D

This table displays the 75 4-dimensional compact euclidean space forms.

Under HOLONOMY there are 4 entries for each space. The first 3 name the arithmetical, geometrical, and isomorphism classes of the holonomy group in the same notations as the preceding tables. The 4th entry designates a representative element of $H^2(F, Z^4)$ in the arbitrary notation used in our data tapes.

Under FROBENIUS MAP is given the fundamental group. Columns A,B,C contain the translation parts to be paired with the corresponding generating matrices of Table C to form a set of generators over Z^4 of the deck group. Under DENOM is the common denominator by which all the vectors under A,B,C are understood to be divided.

Under $H_1(M, Z)$ is given the first homology group of the space. The sequence of integers $k \ l \ m \ \dots$ denotes $(Z/kZ) \oplus (Z/lZ) \oplus (Z/mZ) \ \dots$. In particular, 0 denotes Z .

TABLE D P. 1

HOLONOMY	FROBENIUS MAP				H1(M,Z)
	A	B	C	DENOM	
AC 0	0				0 0 0 0
GC0	0				
C1	0				
1	0			/1	
AC 1	0				2 2 0 0
GC1	0				
C2	1				
1	0			/2	
AC 2	0				2 0 0
GC1	0				
C2	1				
1	0			/2	
AC 7	0	1			4 4 0
GC2	1	0			
C2XC2	1	0			
11	0	0		/2	
AC 7	0	1			2 2 2 0
GC2	0	0			
C2XC2	0	0			
12	1	0		/2	
AC 7	0	1			2 4 0
GC2	1	0			
C2XC2	0	0			
13	1	0		/2	
AC 7	0	1			2 4 0
GC2	1	0			
C2XC2	1	0			
15	1	0		/2	

TABLE D P. 2

HOLONOMY	FROBENIUS MAP				H1(M,Z)
	A	B	C	DENOM	
AC 8	0	1			2 4 0
GC2	0	1			
C2XC2	1	0			
7	0	0		/2	
AC 10	0	1			2 2 0
GC2	0	1			
C2XC2	1	0			
7	0	0		/2	
AC 14	0	1			2 2 0
GC2	0	1			
C2XC2	1	0			
7	0	0		/2	
AC 23	0	0			2 2 0
GC2	0	0			
C2XC2	1	0			
5	0	1		/2	
AC 24	1	0			4 0
GC2	0	0			
C2XC2	1	0			
3	0	1		/2	
AC 27	0				2 0 0
GC3	0				
C4	1				
1	0			/4	
AC 28	0				0 0
GC3	0				
C4	1				
2	0			/4	

TABLE D P. 3

HOLONOMY	FROBENIUS MAP			HI (M,Z)
	A	B	C DENOM	
AC 30	0	0		2 2 0
GC4	0	0		
D4	1	0		
13	0	2	/4	
AC 30	2	0		4 0
GC4	0	0		
D4	1	0		
14	2	2	/4	
AC 30	2	0		4 0
GC4	0	0		
D4	1	0		
15	0	2	/4	
AC 31	2	0		2 0
GC4	0	0		
D4	1	0		
7	2	2	/4	
AC 36	0			3 0 0
GC5	0			
C3	1			
1	0		/3	
AC 37	0			0 0
GC5	0			
C3	1			
1	0		/3	
AC 40	0			0 0
GC6	0			
C6	1			
1	0		/6	

TABLE D P. 4

HOLONOMY	FROBENIUS MAP				H1(M,Z)
	A	B	C	DENOM	
AC 41	0	0			6 0
GC7	0	0			
D3	2	0			
1	0	3		/6	
AC 45	0	0			2 0
GC7	0	0			
D3	2	0			
1	0	3		/6	
AC 47	0	0			2 0
GC7	0	0			
D3	2	0			
i	0	3		/6	
AC 5i	0	0			2 0
GC8	0	0			
D6	1	0			
7	0	3		/6	
AC 53	3	0			0
GC9	0	0			
T	3	0			
i	2	4		/6	
AC 54	3	4			0
GC9	3	0			
T	0	0			
1	4	2		/6	
AC 66	1				2 0 0 0
GC11	0				
C2	0				
1	0			/2	

TABLE D P. 5

HOLONOMY	FROBENIUS MAP				H1(M,Z)
	A	B	C	DENOM	
AC 67	0				0 0 0
GC11	1				
C2	0				
1	0			/2	
AC 69	0				2 2 2 0
GC12	0				
C2	0				
1	1			/2	
AC 77	0	0			2 2 0 0
GC13	0	1			
C2XC2	1	0			
10	0	0		/2	
AC 77	1	0			4 0 0
GC13	0	1			
C2XC2	1	0			
11	0	0		/2	
AC 77	0	0			2 2 0 0
GC13	0	0			
C2XC2	0	1			
20	1	0		/2	
AC 77	1	0			2 0 0
GC13	0	0			
C2XC2	0	1			
21	1	0		/2	
AC 77	1	0			4 0 0
GC13	0	1			
C2XC2	1	1			
27	0	0		/2	

TABLE D P. 6

HOLONOMY	FROBENIUS MAP			HI(M,Z)
	A	B	C DENOM	
AC 77	1	0		2 0 0
GC13	0	1		
C2XC2	0	1		
29	1	0	/2	
AC 78	0	0		2 0 0
GC13	0	1		
C2XC2	1	0		
5	0	0	/2	
AC 78	1	0		2 2 0 0
GC13	0	0		
C2XC2	1	1		
11	0	0	/2	
AC 78	1	0		2 0 0
GC13	0	1		
C2XC2	1	1		
15	0	0	/2	
AC 80	0	0		2 0 0
GC13	0	1		
C2XC2	1	0		
5	0	0	/2	
AC 81	1	1		0 0
GC13	1	0		
C2XC2	0	0		
3	1	0	/2	
AC 84	0	0		2 0 0
GC13	0	0		
C2XC2	0	1		
6	1	0	/2	

TABLE D P. 7

HOLONOMY	FROBENIUS MAP			HI (M,Z)
	A	B	C DENOM	
AC 88	0	0		2 2 2 0
GC14	0	0		
C2XC2	1	0		
20	0	1	/2	
AC 88	1	0		2 4 0
GC14	0	0		
C2XC2	1	0		
21	0	1	/2	
AC 88	1	0		2 2 2 0
GC14	0	0		
C2XC2	0	0		
25	1	1	/2	
AC 88	1	0		2 4 0
GC14	0	0		
C2XC2	1	0		
29	1	1	/2	
AC 96	0	1		2 2 0
GC14	0	0		
C2XC2	0	0		
7	1	1	/2	
AC100	1	0		2 4 4
GC15	0	0		
C2XC2	1	0		
13	0	1	/2	
AC112	0	1	0	2 4 0
GC16	1	0	0	
C2XC2XC2	1	0	0	
278	0	0	1 /2	

TABLE D P. 8

HOLONOMY	FROBENIUS MAP				H1(M,Z)
	A	B	C	DENOM	
AC112	0	1	0		2 2 2 0
GC16	1	0	0		
C2XC2XC2	0	0	0		
282	1	0	1	/2	
AC112	0	1	0		2 4 0
GC16	1	0	0		
C2XC2XC2	1	0	0		
286	1	0	1	/2	
AC112	1	1	0		2 2 0
GC16	0	1	0		
C2XC2XC2	1	0	0		
311	0	0	1	/2	
AC112	0	1	0		2 2 2 0
GC16	1	1	0		
C2XC2XC2	0	0	0		
312	1	0	1	/2	
AC112	0	1	0		2 4 0
GC16	1	1	0		
C2XC2XC2	1	0	0		
316	1	0	1	/2	
AC112	1	1	0		2 2 0
GC16	1	1	0		
C2XC2XC2	1	0	0		
317	1	0	1	/2	
AC112	1	1	0		2 2 0
GC16	0	1	0		
C2XC2XC2	1	0	0		
319	1	0	1	/2	

TABLE D P. 9

HOLONOMY	FROBENIUS MAP				H1(M,Z)
	A	B	C	DENOM	
AC112	0	1	0		2 2 0
GC16	1	1	0		
C2XC2XC2	0	1	0		
376	1	0	1	/2	
AC112	1	1	0		2 2 0
GC16	0	1	0		
C2XC2XC2	0	1	0		
379	1	0	1	/2	
AC130	1	1	0		2 2 4
GC17	0	1	0		
C2XC2XC2	1	0	0		
183	0	0	1	/2	
AC130	1	1	0		2 2 4
GC17	1	1	0		
C2XC2XC2	1	0	0		
189	1	0	1	/2	
AC148	0				2 2 0
GC18	0				
C4	0				
1	1			/4	
AC150	0				2 2 0
GC18	0				
C4	0				
1	1			/2	
AC151	0				4 0
GC18	0				
C4	0				
1	1			/2	

TABLE D P.10

HOLONOMY	FROBENIUS MAP				HI(M,Z)
	A	B	C	DENOM	
AC152	0				4 0
GC18	0				
C4	0				
1	1			/4	
AC156	0	0			2 2 0
GC19	2	0			
C4XC2	0	0			
13	1	2		/4	
AC156	0	0			2 2 0
GC19	2	0			
C4XC2	2	0			
15	1	2		/4	
AC160	0	0			2 0
GC19	0	0			
C4XC2	2	0			
7	1	2		/4	
AC170	0	0			4 0
GC22	1	0			
D4	0	0			
5	1	1		/2	
AC171	0	0			2 0
GC22	0	0			
D4	1	0			
5	1	1		/2	
AC176	2	0			4 4
GC23	0	0			
D4	1	0			
7	0	2		/4	

TABLE D P.11

HOLONOMY	FROBENIUS MAP				H ₁ (M, Z)
	A	B	C	DENOM	
AC209	0				6 0
GC28	0				
C6	0				
1	1			/6	
AC215	0				2 0
GC29	0				
C6	0				
1	1			/6	
AC217	0				2 0
GC29	0				
C6	0				
1	1			/6	
AC219	0	0			2 0
GC30	0	0			
C6XC2	3	0			
11	2	3		/6	
AC257	3	3	0		0
GC40	3	3	0		
TXC2	0	0	0		
7	2	4	3	/6	

Table E

This table is comprised of listings of the source language of the actual versions of the programs used for the main result.

	IDENT	DIAR	
	ENTRY	DIAK,DIAKL,DIALT	
GOA	LDA,I	0,1	NRD
	SWA	F51C	
	SWA	F50D	
	SWA	FD120B	
	SWA	FD150B	
	SWA	F154	
	LDA	1,1	L(MA)
	SWA	LMA1	
	SWA	F150C	
	SWA	F150G	
	LDA	2,1	L(KK)
	SWA	F200B	
	LDA,I	3,1	NRC
	INA,S	-1	
	SWA	F70	
	SWA	FD86G	
	SWA	F109	
	SWA	F130	
	SWA	F149	
	SWA	F150F	
	SWA	F50C	
	SWA	DLICOP.2	
	SWA	DLICS.2	
	SWA	DLNRC	
	MUA	DIALN	
	ANA	-0	
	ADA	LML	
	SWA	DLROP.1	
	SWA	DLRS.1	
	LDA,I	4,1	NCC
	INA,S	-1	
	LQO	F50C	
	ANQ	-0	
	AQJ,GE	*+2	
	SWA	F50C	
	SWA	DKNCC	
	SWA	DKCOP.4	
	SWA	DKCO1.4	
	SWA	DKCS.4	
	MUA	F50D	
	ANA	-0	
	ADA	LMA1	L(MA(1,NCC))
	SWA	F70A	
	SWA	F120	
	SWA	F150	
	SWA	F156	
	INI	5,1	
	STI	RTRN,1	
SW8	BSS	1	
SW9	BSS	1	
F50A	ENA	0	
	STA	KKM	

	SWA	F154A	
	UJP	F51	
* START	NEXT BLOCK		
F50	LDA	KKM	
F50C	ENQ	**	**=NM-1
	AQJ,GE	F200	
	INA	1	
F50B	STA	KKM	
	SWA	F154A	
	LDA	LMA1	
F50D	INA	**	**=NRD
	SWA	LMA1	
	SWA	F150C	
	SWA	F150G	
* FIND	PIVOT ELEMENT		
F51	ENA	FD70	
	SWA	F51A	
	SWA	F70+1	
	ENQ	0	
LMA1	ENI	**2	**=L(MA(1,KK))
	LDI	KKM,1	
	UJP	F51A	
F51C	INI	**2	**=NRD
	INI	1,1	
F51A	STI	**2	**=FD70 OR FD70A
F51B	LDI	KKM,3	
	UJP	F70+1	
FD70	LDA	**3	**=L(MA(1,J))
	AZJ,EQ	F70	
	STI	IP,3	
	STI	JP,2	
	STI	JP1,1	
	ASG,S	1	
	XOA,S	-0	
	ASG	2	
	UJP	FGI+2	
	SHAQ	-24	
	ENA	FD70A	
	SWA	F70+1	
	SWA	F51A	
	STI	FD70A,2	
	UJP	F70	
FD70A	LDA	**3	
	AZJ,FQ	F70	
	ASG,S	1	
	XOA,S	-0	
	AQJ,GE	F70	
	STI	IP,3	
	STI	JP,2	
	STI	JP1,1	
	ASG	2	
	UJP	FGI+2	
	SHAQ	-24	
F70	YSI	**3	**=NRC-1
	UJP	**	**=FD70 OR FD70A

F70A	ISG	** , 2	** = L (MA (1 , NCC))
	UJP	F51C	
FGI	SHAQ	24	
	AZJ , EQ	F200A	
	STA	ABSIV	
	LDI	JP , 1	
FGI5	STI	FD120A , 1	(1) = L (MA (1 , JP)
	STI	FD86E , 1	
	STI	FGI7 , 1	
FGI6	LDI	IP , 1	(1) = IP - 1
	STI	F85 , 1	
FGI7	INI	** , 1	
	STI	F86 , 1	
	ASG	2	
	UJP	FGJ	
	ENI	F85B , 1	
	ENI	FD109A + 1 , 2	
	ENI	F1129A , 3	
	UJP	FGK	
FGJ	LDA , I	F86	
	AZJ , GE	FGJ1	
	ENI	F86D , 1	
	UJP	** + 2	
FGJ1	ENI	F86C , 1	
	ENI	F120 , 2	
	ENI	F150B , 3	
FGK	STI	F85A , 1	
	STI	FD109A , 2	
	STI	F130A , 3	
SW1	BSS	1	
* IS ROW	0		
FD120	LDI	LMA1 , 2	
	LDI	KKM , 1	
	UJP	FD120A	
FD120B	INI	** , 2	** = NRD
	INI	1 , 1	
FD120A	ISE	** , 2	** = L (MA (1 , JP))
	UJP	** + 2	
	UJP	F120	
F85	LDA	** , 2	** = IP - 1
	-ZJ , EQ	F120	
F85A	UJP	**	** = F85B OR F86C OR F86D
F85B	SHAQ	-24	
F86	DVA	**	** = L (MA (IP , JP))
F86C	XOA , S	-0	
F86D	STA	KWO	
	STI	FD86F , 2	
	LDI	KKM , 3	
FD86E	LDA	** , 3	** = L (MA (1 , JP)
	MUA	KWO	
FD86F	RAO	** , 3	** = L (MA (1 , J)
FD86G	ISI	** , 3	** = NRC - 1
	UJP	FD86E	
SW2	BSS	1	
FD109A	UJP	**	** = ** + 1 OR F120

	LDQ	ARSIV	
	STI	FD100,2	
	LDI	KKM,3	
FD109	LDA	** ,3	**=L(MA(1,J))
	AZJ,EQ	F109	
	ASG,S	1	
	XQA,S	-0	
	AQJ,GE	F109	
	STI	IP,3	
	SHAQ	-24	
F109	ISI	** ,3	**=NRC-1
	UJP	FD109	
	LDA	ARSIV	
	NCC,EQ	F120	
	STI	JP,2	
	STI	JP1,1	
	UJP	FGI	
F120	ISG	** ,2	**=L(MA(1,NCC))
	UJP	FD120B	
* IS COLUMN 0			
FD130	LDA	F86	
	SWA	F121A	
	LDA	JP	
	SWA	F121	
	LDA	IP	
	SWA	**2	
FD130A	LDI	KKM,3	
F121B	ISE	** ,3	
	UJP	F121	
	UJP	F130	
F121	LDA	** ,3	
	AZJ,EQ	F130	
	SHAQ	-24	
F121A	DVA	**	**=L(MA(IP,JP))
	STQ,I	F121	
	STA	KWO	
SW3	BSS	1	
	LDA,I	F121	
	AZJ,EQ	F130	
	ASG,S	1	
	XQA,S	-0	
	STA	ARSIV	
	STI	IP,3	
	UJP	FGI6	
F130	ISI	** ,3	**=NRC-1
	UJP	F121B	
F130A	UJP	**	**=F1129A OR F150B
F1129A	LDA	F121A	
	SWA	F1131A	
	LDA	JP	
	SWA	FD150A	
	SWA	F135B	
	LDA	IP	
	SWA	FD149A	
	SWA	F135A	

* DIVIDES REMAINING ELEMENTS

FD150	LDI	LMA1,2	
	LDI	KKM,1	
	UJP	FD150A	
FD150B	INI	** ,2	**=NRD
	INI	,1	
FD150A	ISE	** ,2	**=L(MA(1,JP))
	UJP	*+2	
	UJP	F150	
	STI	F1131,2	
FD149	LDI	KKM,3	
FD149A	ISE	** ,3	**=IP-1
	UJP	*+2	
	UJP	F149	
F1131	LDA	** ,3	**=L(MA(1,J))
	AZJ, EQ	F149	
	SHAQ	-24	
F1131A	DVA	**	**=L(MA(IP,JP))
	SHAQ	24	
	AZJ, EQ	F149	
	STA, I	F1131	
	ASG, S		
	XOA, S	-0	
	STA	ABSIV	
	STQ	KWO	
	LDA, I	F1131A	
F135A	STA	** ,2	**=IP-1
	MUA	KWO	
	XOA, S	-0	
F135B	STA	** ,3	**=L(MA(1,JP))
	STI	JP,2	
	STI	JP1,1	
SW4	SSS	1	
	STI	IP,3	
SW5	BSS	1	
	LDA	ABSIV	
	UJP	FGI+3	
F149	ISI	** ,3	**=NRC-1
	UJP	FD149A	
F150	ISG	** ,2	**=L(MA(1,NCC))
	UJP	FD150B	
* MOVE ROWS AND COLUMNS			
F150B	ENG	0	
	LDA	JP	
	SWA	F150E	
	LDI	KKM,3	
F150C	LDA	** ,3	**=L(MA(1,KK))
F150D	STQ, I	*-1	
F150E	STL	** ,3	**=L(MA(1,JP))
F150F	ISI	** ,3	**=NRC-1
	UJP	F150C	
F150G	ENI	** ,2	
	LDI	IP,1	
	STI	F155,1	
	UJP	*+2	

F154	INI	** , 2	**=NRD
F154A	LDA	** , 2	**=KK-1
	STQ, I	*-1	
F155	STA	** , 2	**IP-1
F156	ISG	** , 2	**=L(MA(1,NCC))
	UJP	F154	
	LDA	ABSIV	
	LDI	KKM, 2	
	STA, I	LMA1	
SW6	BSS	1	
SW7	BSS	1	
	UJP	F50	
F200	INA	1	
	UJP	**+2	
F200A	LDA	KKM	
F200B	STA	**	
RTRN	UJP	**	
KWO	BSS	1	
ARSIV	BSS	1	
IP	OCT	0	
JP	OCT	0	
KKM	OCT	0	
DIALI	UJP	**	
	LDI	*-1, 1	
	LDA	DNQP	
	STA	SW2	
	STA	SW5	
	STA	SW6	
	STA	SW8	
	LDA	DLISW1	
	STA	SW1	
	LDA	DLISW3	
	STA	SW3	
	STA	SW4	
	LDA	DLISW7	
	STA	SW7	
	LDA	DLSW9	
	STA	SW9	
	LDA, I	0, 1	NRDL
	SWA	DIALN	
	LDA	1, 1	
	SWA	LML	
	INI	2, 1	
	UJP	GOA	
DIACL	UJP	**	
	LDI	*-1, 1	
	LDA	DLSW3	
	STA	SW3	
	STA	SW4	
	LDA	DLSW7	
	STA	SW7	
	LDA	DLSW9	
	STA	SW9	
	LDA, I	0, 1	NRDL
	SWA	DIALN	

	SWA	DLROP.2	
	SWA	DLRS.2	
	LDA	1,1	
	SWA	LML	
	INI	2,1	
DIKAK	UJP	DIKAK.1	
	UJP	**	
	LDI	*-1,1	
	LDA	DNOP	
	STA	SW3	
	STA	SW4	
	STA	SW7	
	STA	SW9	
DIKAK.1	LDA	DKSW1	
	STA	SW1	
	LDA	DKSW2	
	STA	SW2	
	LDA	DKSW5	
	STA	SW5	
	LDA	DKSW6	
	STA	SW6	
	LDA	DKSW8	
	STA	SW8	
	LDA, I	0,1	NRDK
	SWA	DIKAKIN	
	LDA	1,1	L(MK)
	SWA	LMK	
	INI	2,1	
	UJP	GOA	
DIKAKIN	ENA	**	**=NRDK
	SWA	IDM6	
DKNCC	ENA	**	**=NCC-1
	SWA	IDM9	
	SWA	IDM15	
LMK	ENI	** , 1	**=L(MK)
	RTJ	IDM	
	UJP	SW9	
DIALN	ENA	**	**=NRDL
	SWA	IDM6	
DLNRC	ENA	**	**=NRC-1
	SWA	IDM9	
	SWA	IDM15	
LML	ENI	** , 1	**=L(ML)
	RTJ	IDM	
	UJP	F50A	
IDM	UJP	**	
	ENI	0,3	
	ENI	0,2	
	ENA	0	
	ENQ	1	
	UJP	*-2	
IDM6	INI	** , 1	
	STI	*+1,1	
	STA	** , 3	
IDM9	ISI	** , 3	

	UJP	*-2	
	STI	**+1,1	
	STQ	** ,2	
IDM15	ISI	** ,2	
	UJP	IDM6	
	UJP ,I	IDM	
JPSET	LDA	JPI	
	MUA	DIAKIN	
	ANA	-0	
	ADA	LMK	
	SWA	DKCOP.1	L(MK(1,JP))
	SWA	DKCO1.1	
	SWA	DKCS.1	
	UJP	SW1+1	
IPSET	LDA	IP	
	MUA	DIALN	
	ANA	-0	
	ADA	LML	L(ML(1,IP))
	SWA	DLICOP.1	
	SWA	DLICS.1	
	UJP	SW1+1	
DKCOP	TIA	1	
	MUA	DIAKIN	
	ANA	-0	
	ADA	LMK	
	SWA	**+4	
	ENI	0,3	
DKCOP.1	LDA	** ,3	**=L(MK(1,JP))
	MUA	KWO	
	RAD	** ,3	
DKCOP.4	ISI	** ,3	**=NCC-1
	UJP	*-4	
	UJP	SW2+1	
DKCO1	TIA	1	
	MUA	DIAKIN	
	ANA	-0	
	ADA	LMK	
	SWA	**+3	
	ENI	0,3	
DKCO1.1	LDA	** ,3	**=L(MK(1,JP))
	RAD	** ,3	
DKCO1.4	ISI	** ,3	**=NCC-1
	UJP	*-3	
	UJP	SW5+1	
DKCS	LDA	KKM	
	MUA	DIAKIN	
	ANA	-0	
	ADA	LMK	
	SWA	**+2	
	ENI	0,3	
	LDA	** ,3	
DKCS.1	LDO	** ,3	**=L(MK(1,JP))
	STO ,I	*-2	
	STA ,I	*-2	
DKCS.4	ISI	** ,3	**=NCC=1

	UJP	*-5	
	UJP	SW7	
DLROP	UJP	**	
	LDA	IP	
	SWA	*+5	
	STI	*+6,3	
	LDI	LML,3	
	UJP	*+2	
DLROP.2	INI	** ,3	**=NRDL
	LCA	** ,3	**=IP-1
	MUA	KW0	
	RAD	** ,3	**=I-1
DLROP.1	ISG	** ,3	**=L (ML (1,NRC))
	UJP	*-5	
	LDI	*-3,3	
	UJP,I	DLROP	
DLRS	LDA	IP	
	SWA	*+7	
	LDA	KKM	
	SWA	*+4	
	LDI	LML,3	
	UJP	*+2	
DLRS.2	INI	** ,3	**=NRDL
	LDA	** ,3	
	LDQ	** ,3	
	STQ,I	*-2	
	STA,I	*-2	
DLRS.1	ISG	** ,3	**=L (ML (1,NRC))
	UJP	*-6	
	UJP	SW7+1	
DLICOP	UJP	**	
	STI	DLICOP.3,3	
	TIA	3	
	MUA	DIALN	
	ANA	-0	
	ADA	LML	
	SWA	*+2	
	ENI	0,3	
	LDA	** ,3	
	MUA	KW0	
DLICOP.1	RAD	** ,3	**=L (ML (1,IP))
DLICOP.2	ISI	** ,3	**=NRC-1
	UJP	*-4	
DLICOP.3	ENI	** ,3	
	UJP,I	DLICOP	
DLICS	LDA	KKM	
	MUA	DIALN	
	ANA	-0	
	ADA	LML	
	SWA	*+2	
	ENI	0,3	
	LDA	** ,3	
DLICS.1	LDQ	** ,3	**=L (ML (1,IP))
	STQ,I	*-2	
	STA,I	*-2	

```
DLICS.2 ISI          ** ,3          **=NRC-1
      UJP          *-5
      UJP          SW7+1
DNOP   ISI          0
DKSW1  UJP          JPSET
DKSW2  UJP          DKCOP
      SW5  UJP          DKCO1
DKSW6  UJP          DKCS
DKSW8  UJP          DIAKTN
DLSW3  RTJ          DLR00
DLSW7  UJP          DLRS
DLSW9  UJP          DIALN
DLISW1 UJP          IPSET
DLISW3 RTJ          DLICOP
DLISW7 UJP          DLICS
      JPI   OCT          0
      END
```

	IDENT	MMPYR	
	ENTRY	MMPY, MMPY1, MMPY2, MMPZ, MMPA	
MMPY1	UJP	**	CALL MMPY1(IA, IB, IC, I1, J1, K1,
	LDI	*-1,1	OR
	LDA, I	0,1	DITTO ..K1,0,A,B,C)
	SWA	IA	
	LDA, I	1,1	
	SWA	IB	
	LDA, I	2,1	
	SWA	IC	
MM2	LDA, I	3,1	
	INA, S	-1	
	SWA	I1	
	LDA, I	4,1	
	INA, S	-1	
	SWA	J1	
	LDA, I	5,1	
	SWA	K1	
	LDA, I	6,1	
	AZU, EQ	**+2	
	UJP	7,1	
	INI	7,1	
	UJP	MM3	
MMPY2	UJP	**	CALL MMPY2(I1, J1, K1, 1)
	LDI	*-1,1	OR
	INI	-3,1	
	UJP	MM2	CALL MMPY2(I1, J1, K1, 0, A, B, C)
MMPY	UJP	**	
	LDI	*-1,1	CALL MMPY(A, B, C)
MM3	LDA	0,1	
	SWA	A	
MM4	LDA	1,1	
	SWA	JB	
	LDA	2,1	
	SWA	JC	
	INI	3,1	
	STI	RTRN, 1	
	ENA	1	
	STA	J	
	UJP	A.	
	INA	1	
	STA	J	
IB	ENA	**	
	RAD	JB	
IC	ENA	**	
	RAD	JC	
A.	INI	0,1	
A	ENA	**	
	SWA	A1	
	UJP	AAA	
AA	ENA	1	
	RAD	A1	
AAA	ENA	0	
	STA	X1	
	INI	0,3	

	ENI	0,2	
	UJP	A1	
IA	INI	** ,2	
A1	LDA	** ,2	
JR	MUA	** ,3	
	RAD	X1	
J1	ISI	** ,3	
	UJP	IA	
	LDA	X1	
JC	STA	** ,1	
I1	ISI	** ,1	
	UJP	AA	
	LDA	J	
K1	ASG	**	
	UJP	IB-2	
RTRN	UJP	**	
J	BSS	i	
X1	BSS	l	
MMPA	UJP	**	CALL MMPA(A)
	LDI	*-1,1	
	LDA	0,1	
	SWA	A	
	UJP	1,1	
MMPZ	UJP	**	CALL MMPZ(B,C)
	LDI	*-1,1	
	INI	-1,1	
	UJP	MM4	
	END		

	IDENT ENTRY	IGCDR IGCD
	UJP	**
	LDI	*-1,1
	LDA,I	0,1
	AZJ,EQ	Z
	ASG,S	0
	XOA,S	-0
	STA	M
	ENA	0
	LDD,I	1,1
	OSG,S	0
	XOQ,S	-0
0	DVA	M
	QSE,S	0
	UJP	L
	LDA	M
	UJP	2,1
L	LDA	M
	STQ	M
	SHAQ	-24
	UJP	D
Z	L,I	1,1
	ASG,S	0
	XOA,S	-0
	UJP	2,1
M	BSS	1
	END	

	IDENT ENTRY ENTRY ENTRY COMMON	POLYR CLSTK,FREPOL PML1,CKPHD1 PDCOD HOMOG,UNF 5	
STACK	BSS	1005	
	PRG		
D24	DEC		
ADTMS	LCA	0,1	
ADTMA	LDA	2,1	
MCTP	UJP	MID+1	
MNCTP	IJI	0	
PT1	OCT	0	
PT2	OCT	0	
TEMP	BSS	10	
PRTRN	UJP	**	
CLSTK	UJP	**	
	ENA	STACK+1	
	STA	STACK	
CLSTK1	TAI	1	
	INA	3	
	STA	0,1	
	ISG	STACK+999,1	END OF STACK
	UJP	CLSTK1	
	ENA	0	
	STA	0,1	
	UJP,I	CLSTK	
NXCEL	UJP	**	
	LDD,I	STACK	
	LDA	STACK	
	ASE	0	
	UJP	**2	
	UJP	STKOV	
	ANQ	77777B	
	STQ	STACK	
	UJP,I	NXCEL	
FRECEL	UJP	**	
	SWA	**2	
	LDA	STACK	
	SWA	**	
	LDA	*-1	
	SWA	STACK	
	UJP,I	FRECEL	
FREPOL	UJP	**	
	LDA,I	*-1	
	SWA	*+1	
	LDA	**	
	ANA	77777B	
	AZJ,EQ	FREPOL1	
	LDA	STACK,2	
	SWA	STACK	
	TAI	1	
	LDA	0,1	
	ASE	0	
	UJP	*-3	

	TIA	2	
	SWA	0,1	
FREPOL1	LDI	FREPOL,1	
	UJP	1,1	
CRPHD1	UJP	**	CALL CRPHD1(N,MV,MC,P)
	LDI	**1,1	
	LDA	3,1	
	SWA	P	
	LDA	2,1	
	SWA	MC	
	LDA	1,1	
	SWA	MV	
	LDA,1	0,1	
	INA,S	-1	
	SWA	N	
	INI	4,1	
	STI	PRTRN,1	
	ENI	0,3	
	ENI	0,1	
CRP1	RTU	NOCEL	
	TAI	2	
MC	LDA	**1	
	STA	2,2	
MV	LDA	**1	
	STA	1,2	
	STI	**1,2	
	STI	**3	
	TIA	2	
	TAI	3	
N	ISI	**1	
	UJP	CRP1	
P	STI	**2	
	UJP,1	PRTRN	
AD12	UJP	**	(1)=LOC(1 TERM P)
	ENA	0	
	STA	TEMP	
AD12A	ISE	0,1	
	UJP	**2	
	UJP	ADRTRN	
	LDQ	1,1	
AD12B	ENI	**3	**=LOC(EXTL LINK VAR TO Q)
AD12D	LDA	0,3	
	TAI	2	
	ISE	0,2	
	UJP	**2	
	UJP	ADNM	
	LDA	1,2	
	AQU,EQ	ADTM	
	STI	**1,2	
	ENI	**3	
	UJP	AD12D	
ADTM	BSS	1	LDA OR LCA 2,1
	ADA	2,2	
	AZU,EQ	ADTZ	
	STA	2,2	

ADTZ	UJP LDA SWA TIA RTJ UJP	ADNPT 0,2 0,3 2 FRECEL ADNPT	
ADNM	RTJ TAI LDA SWA LDA STA	NXCEL 2 TEMP 0,2 1,1 1,2	
ADNMI	SSS STA STI	1 2,2 TEMP,2	LDA OR LCA 2,1
ADNPT	LDA TAI UJP	0,1 1 AD12A	
ADRTRN	LDA QSE UJP UJP LDI LDA ANA AZU, EQ TAI UJP SHAQ SWA UJP, I	TEMP 0 *+2 AD12 AD12B,2 0,2 77777B *+3 2 *-4 24 0,2 AD12 ** *-1,1 3,1 *+3 DIMS *+2 ADTMA ADTM ADNMI 2,1 AD12B 1,1 MI 0,1 M,8 4,1 PRTRN,1	
PML1	UJP LDI LDA, I AZU, GE LDA UJP LDA STA STA LDA SWA LDA, I SWA LDA, I SWA INI STI	** *-1,1 3,1 *+3 DIMS *+2 ADTMA ADTM ADNMI 2,1 AD12B 1,1 MI 0,1 M,8 4,1 PRTRN,1	CALL PML1(P,Q,R,ISW)
M,8	ENA	**	**=LOC(1 TERM P)
M,9	ASE UJP UJP, I STA INA SWA	0 *+2 PRTRN PT2 1 MIAP	(PT2)=LOC(CURRENT TERM P)

	INA	1	
	SWA	M1H	
	ENA	0	
	STA	PT1	
	LDA, I	M1AP	
	AZJ, NE	M.99	
	LDA	MCTP	
	STA	M1APCT	
	UJP	M1	
M.99	LDA	MNCTP	
	STA	M1APCT	
M1	ENI	**2	**=LOC(1 TERM Q)
M1.1	ISE	0.2	
	UJP	**2	(2)=LOC(CURRENT TERM Q)
	UJP	M2	
	LDA	1.2	
	STA	TEMP	
	ENI	0.3	SHIFT COUNTER
M1APCT	BSS	1	
MIA	ENA	0	
	LDA	TEMP	
	QSE, S	0	
	UJP	**2	
	UJP	M1F	
	DVA	D24	
	STA	TEMP	
M1AP	LDA	**	**=LOC(VF TERM P)
	AQJ, GE	M1B	
	ENI	1.3	
	STQ	TEMP, 3	
	UJP	M1A	
M1B	STQ	TEMP +9	
	LDA	TEMP	
	MUA	D24	
	ADA	TEMP +9	
	MUA	D24	
	ADA, I	M1AP	
M1E	ISE	0.3	
	UJP	M1C	
	UJP	M1D	
M1C	MUA	D24	
	ADA	TEMP, 3	
	ISD	1.3	
	UJP	M1C	
M1D	STA	TEMP	
	RTJ	MXCEL	
	SWA	**1	
	ENI	**3	
	LDA	PT1	
	SWA	0.3	
	STI	PT1, 3	
	LDA	TEMP	
	STA	1.3	
M1H	LDA	**	P COEF FIELD
	MUA	2.2	

	STA	2,3
	LDA	0,2
	TAI	2
	UJP	M1,1
MIF	LDA,I	M1AD
	UJP	MIF
ME	LDI	PT1,1
	RTJ	AD12
	RTJ	FREPOL
	77	PT1
	LDA,I	PT2
	UJP	M19
PDCOD	UJP	**
	LDI	**1,1
	LDA,I	4,1
	SWA	PDNTM
	LDA	3,1
	SWA	PDNTI
	LDA	2,1
	SWA	PDMC
	LDA	1,1
	SWA	PDMV
	END	0,2
	LDA,I	0,1
PLP	TAI	3
	ISE	0,3
	UJP	**2
	UJP	PDNT
	LDA	1,3
PDMV	STA	**2
	LDA	2,3
PDMC	STA	**2
	LDA	0,3
	INI	1,2
PDNTM	ISG	**2
	UJP	PLP
PDNT	TIA	2
PDNTI	STA	**
	UJP	5,1
HOMOG	UJP	**
	LDI	**1,1
	LDA	0,1
	SWA	**2
	SWA	**8
	LDI	**2
	ISG	1,2
	UJP	1,1
	LDA	1,2
	ASE,S	0
	UJP	1,1
	LDA	0,2
	SWA	**
	TIA	2
	RTJ	FRECEL
	UJP	1,1

CALL PDCOD(MP,MV,MC,NT,NTM)

CALL HOMOG(MP)

```

UNHMG  UJP      **                CALL UNHMG(MP,MC)
        LDI     *+1,1
        RTJ     NXCEL
        TAI     2
        LDA,1   1,1
        STA     2,2
        ENA     0
        STA     1,2
        LDA     0,1
        SWA     *+1
UNHMGCP LDA     **
        SWA     0,2
        STI,1   UNHMGCP,2
        UJP     2,1
        RTJ     ABNORMAL
        UJP     **
        EXT     ABNORMAL
        END

```

```

FUNCTION IDET4(M)
DIMENSION M(4,4),M1(3,3)
DO 10 I=1,3
DO 10 J=1,3
10 M1(I,J)=M(I+1,J+1)
ISW=1
IDET4=0
DO 100 I=1,4
IF(M(I).EQ.0)GO TO 15
IDET=IDETS3(M1)
CALL PML1(M(I),IDET,IDET4,ISW)
CALL FREPOL(IDET)
15 IF(ISW.GT.0)GO TO 20
ISW=1
GO TO 25
20 ISW=-1
25 IF(I.EQ.4)RETURN
DO 100 J=1,3
100 M1(I,J)=M(I,J+1)
END
FUNCTION IDETS3(M)
DIMENSION M(3,3),M1(2,2)
DO 10 I=1,2
DO 10 J=1,2
10 M1(I,J)=M(I+1,J+1)
ISW=1
IDETS3=0
DO 100 I=1,3
IF(M(I).EQ.0)GO TO 15
IDET=IDETS2(M1)
CALL PML1(M(I),IDET,IDETS3,ISW)
CALL FREPOL(IDET)
15 IF(ISW.GT.0)GO TO 20
ISW=1
GO TO 25
20 ISW=-1
25 IF(I.EQ.3)RETURN
DO 100 J=1,2
100 M1(I,J)=M(I,J+1)
END
FUNCTION IDETS2(M)
DIMENSION M(2,2)
IDETS2=0
CALL PML1(M(1,1),M(2,2),IDETS2,1)
CALL PML1(M(1,2),M(2,1),IDETS2,-1)
RETURN
END

```

```
SUBROUTINE NITTY(NA,NB,MU,MA,MB,MDEL)
DIMENSION MA(5,5,3),MB(5,5,3),MDEL(75,25)
NANB=NA*NB
NANSMU=NANB*MU
DO 170 I1=1,NANBMU
DO 170 J1=1,NANB
170 MDEL(I1,J1)=0
DO 200 INU=1,MU
I0=(INU-1)*NANB
DO 200 I1=1,NB
J2=(I1-1)*NA
I2=J2+I0
DO 200 I3=1,NA
I4=I3+I2
DO 180 J3=1,NA
J4=J2+J3
180 MDEL(I4,J4)=MA(I3,J3,INU)
DO 200 J1=1,NB
J4=(J1-1)*NA+I3
200 MDEL(I4,J4)=MDEL(I4,J4)-MB(J1,I1,INU)
RETURN
END
```

```

SUBROUTINE MATP(ND,NP,MA,MB)
DIMENSION MA(5,5),MB(5,5),MI(5,5)
JSW=1
GO TO 10
ENTRY MATI
JSW=2
10 NP1=NP
IF(NP1)15,15,20
15 DO 18 I=1,ND
DO 16 J=1,ND
16 MB(I,J)=0
18 MB(I,I)=1
RETURN
20 DO 25 I=1,ND
DO 23 J=1,ND
23 MB(I,J)=MA(I,J)
GO TO(25,24),JSW
24 MB(I,I)=MB(I,I)+1
25 CONTINUE
NP1=NP1-1
IF(NP1)80,80,30
30 CALL MMPA(MA)
40 CALL MMPZ(MB,MI)
JSW=1
GO TO(60,42),JSW
42 DO 43 I=1,ND
43 MI(I,I)=MI(I,I)+1
GO TO 60
50 CALL MMPZ(MI,MB)
JSW=2
GO TO(60,52),JSW
52 DO 53 I=1,ND
53 MB(I,I)=MB(I,I)+1
60 NP1=NP1-1
IF(NP1)70,70,65
65 GO TO(50,40),JSW
70 GO TO(75,80),JSW
75 DO 76 I=1,ND
DO 76 J=1,ND
76 MB(I,J)=MI(I,J)
80 RETURN
END

```

```
SUBROUTINE WORDC(ND,LI,LE,MG,MGA)
DIMENSION MG(5,5,3),MGA(5,5),MI(5,5),M2(5,5),LI(4),LE(4)
LI=LI(1)
CALL MATP(ND,LE(1),MG(1,1,LI),MGA)
ISW=1
DO 10 I=2,4
LI=LI(I)
IF(LI.EQ.0)GO TO 20
CALL MATP(ND,LE(I),MG(1,1,LI),MI)
GO TO(5,6),ISW
5 CALL MRPY(MGA,MI,M2)
ISW=2
GO TO 10
6 CALL MRPY(M2,MI,MGA)
ISW=1
10 CONTINUE
20 IF(ISW.EQ.1)RETURN
CALL MATP(ND,1,M2,MGA)
RETURN
END
```

```

SUBROUTINE WORDV(LI,LE,MX,M1)
COMMON MA(5,5,5),ND,NDMU
DIMENSION LI(4),LE(4),MX(5,25),M1(5,5)
EQUIVALENCE (MA,MCOM(2956)),(ND,MCOM(3911)),(NDMU,MCOM(3912))
DIMENSION M2(5,5)
MWL=4
DO 505 I=1,ND
DO 505 J=1,NDMU
505 MX(I,J)=0
K1=1
K2=LI(K1)
K3=LE(K1)
J1=ND*(K2-1)
DO 510 J2=1,ND
J3=J1+J2
DO 510 I2=1,ND
IF(I2-J2)509,508,509
508 MX(I2,J3)=1
GO TO 510
509 MX(I2,J3)=0
510 M1(I2,J2)=MA(I2,J2,K2)
ISW=1
GO TO 540
520 GO TO(521,526),ISW
521 DO 525 J2=1,ND
J3=J1+J2
DO 525 I2=1,ND
525 MX(I2,J3)=MX(I2,J3)+M1(I2,J2)
CALL MMPY(M1,MA(I,I,K2),M2)
ISW=2
GO TO 540
526 DO 530 J2=1,ND
J3=J1+J2
DO 530 I2=1,ND
530 MX(I2,J3)=MX(I2,J3)+M2(I2,J2)
CALL MMPY(M2,MA(I,I,K2),M1)
ISW=1
540 K3=K3-1
IF(K3)545,545,520
545 IF(K1-MWL)550,550,500
550 K1=K1+1
K2=LI(K1)
IF(K2)600,600,555
555 J1=ND*(K2-1)
K3=LE(K1)
GO TO 520
600 GO TO(610,605),ISW
605 DO 606 J=1,ND
DO 606 I=1,ND
606 M1(I,J)=M2(I,J)
610 RETURN
END
FINIS

```

```

PROGRAM VB
COMMON MA(5,5,3),MB(2,2,3),MDEL(18,6),MLI(18,18),ISOC(2),MCC(6),
C MCX(6),LAI(4,3,8),LAE(4,3,8)
1 FORMAT(7I3,2A4)
20 FORMAT(6H1GEOCL, I3,5X,5HISOCL, 1X,2A4,5X,4HNDEC, I3/5H GENS, I3,5X,
1 2HNA, I3,5X,2HNB, I3,5X,4HNAUT, I3)
3 FORMAT(5H0IDEC, I3,5X,4HCHRK, I3,5X,5HINVAR,6I3)
4 FORMAT(4H0IAC, I3)
5 FORMAT(1X/6H*GENER, I3)
6 FORMAT(12X,8I3)
7 FORMAT(26I3)
8 FORMAT(1X/4H*AUT, I3)
CALL DIAGLI1(18,18,MDEL, KK, MLI)
MWL=4
CALL MMPY1(18,6,5,0,0,0,1)
CALL MMPA(MLI)
NZERO=0
NEGI=-1
N888=888
READ 7, IAC
C GEN: GEO: CLASS
20 READ 1, IFL, NGC, NDEC, MU, NA, NB, NAUT, ISOC
CALL SSWICH(1, IP)
GO TO(21,22), IP
21 PUNCH1, IFL, NGC, NDEC, MU, NA, NB, NAUT, ISOC
22 IF(IFL.EQ.999) GO TO 25
WRITE TAPE 11, (NZERO, I=1,10)
REWIND 11
STOP
25 NANS=NA*NB
NANSMU=NANS*MU
NAI=NA+1
ND=NA+NB
DO 26 K=1, NAUT
DO 26 J=1, MU
26 READ 7, (LAI(I, J, K), LAE(I, J, K), I=1, MWL)
WRITE TAPE 11, IFL, NGC, ISOC, NDEC, MU, NA, NB, ND, NAUT,
1 (LAI(I, J, K), LAE(I, J, K), I=1, MWL), J=1, MU), K=1, NAUT)
PRINT 2, NGC, ISOC, NDEC, MU, NA, NB, NAUT
DO 28 K=1, NAUT
PRINT 6, R
DO 28 J=1, MU
28 PRINT 6, (LAI(I, J, K), LAE(I, J, K), I=1, MWL)
DO 30 K=1, MU
DO 30 I=NAI, ND
DO 30 J=1, NA
30 MA(I, J, K)=0
DO 300 IDEC=1, NDEC
DO 60 K=1, MU
READ 7, ((MA(I, J, K), J=1, NA), I=1, NA)
60 READ 7, ((MB(I, J, K), J=1, NB), I=1, NB)
CALL REV. (NA, ND, MU, MA, MB, MDEL)
CALL DIAGLI(NANSMU, NANS)
PRINT3, IDEC, KK, (MDEL(I, I), I=1, KK)
DO 70 K=1, MU

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```

DO 70 I=1,NS
  I1=I+NA
DO 70 J=1,NS
  J1=J+NA
70 MA(I1,J1,K)=MB(I,J,K)
  CALL HMPY2(NANSMU, KK, 1, 1)
DO 105 I=1, KK
105 MCC(I)=0
110 CALL HMPZ(MCC, MCX)
  I0=0
DO 120 K=1, MU
DO 120 J=NA1, ND
DO 120 I=1, NA
  I0=I0+1
120 MA(I, J, K)=MCX(I0)
  IAC=IAC+1
  PRINT 4, IAC
  GO TO(121, 122), IP
121 PUNCH 7, N888, IAC, ND
122 DO 125 K=1, MU
  GO TO(123, 124), IP
123 PUNCH 7, ((MA(I, J, K), J=1, ND), I=1, ND)
124 PRINT 5, K
  DO 125 I=1, ND
125 PRINT 6, (MA(I, J, K), J=1, ND)
  WRITE TAPE 11, IDEC, IAC, (((MA(I, J, K), I=1, ND), J=1, ND), K=1, MU)
210 KINC=1
215 KCI=MCC(KINC)+1
  IF(MCICLT, MDEL(KINC, KINC)) GO TO 225
  KINC=KINC+1
  IF(KINC, LE, KK) GO TO 215
  GO TO 300
225 KIN=KINC-1
DO 235 I=1, KIN
235 MCC(I)=0
  MCC(KINC)=MCI
  GO TO 110
300 CONTINUE
  WRITE TAPE 11, NEG1, NEG1, (((NEG1, I=1, ND), J=1, ND), K=1, MU)
  GO TO 20
END

```

```

PROGRAM VC
COMMON MX(75,25),MK(25,25),ISOC(2),MG(5,5,3),MH(5,5,3),MV(100),
1 MC(100),M4(4,4),MP(5),MR,KT,MGA(5,5,24),LAI(4,3,8),LAE(4,3,8)
2 FORMAT(6H0GEOCL,13,5X,5HISOCL,1X,2A4,5X,4HGENS,13,5X,3HDIM,13)
3 FORMAT(6H0POSS EQU,SI4/10X,4HCOEF,5X,3HVAR)
4 FORMAT(10X,I6,5X,SI4)
5 FORMAT(6H0OVFL0,4I4)
CALL DIAGKI(75,25,MX,KT,MK)
CALL MNPY1(5,5,5,0,0,0,1)
MNL=4
20 READ TAPE 11,IFL,NGC,ISOC,NDEC,MU,NA,NB,ND,NAUT,
1 ((LAI(I,J,K),LAE(I,J,K),I=1,MNL),J=1,MU),K=1,NAUT)
IF(IFL.EQ.999)GO TO 25
REWIND 11
STOP
25 PRINT 1,NGC,ISOC,MU,ND
ND2=ND**2
MXR=MU*ND2
30 READ TAPE 11,IDEC,IAC,(((MG(I,J,K),I=1,ND),J=1,ND),K=1,MU)
IF(IAC.LE.0)GO TO 20
CALL MNPY2(ND,ND,ND,1)
DO 230 K=1,NAUT
KI=(K-1)*MU
DO 230 J=1,MU
J1=KI+J
230 CALL WORDC(ND,LAI(1,J,K),LAE(1,J,K),MG,MGA(1,1,J1))
KBS=0
35 KBS=KBS+1
READ TAPE 11,JDEC,JAC,(((MH(I,J,K),I=1,ND),J=1,ND),K=1,MU)
IF(JAC.GT.0)GO TO 100
IF(KSS.EQ.1)GO TO 20
DO 40 I=1,KBS
40 BACKSPACE 11
GO TO 30
100 DO 400 KA=1,NAUT
KAI=(KA-1)*MU+1
CALL NETTY(ND,ND,MU,MGA(1,1,KAI),MH,MX)
CALL DIAGK(MXR,ND2)
IF(KK.EQ.ND2)GO TO 400
MC1=KK+1
CALL CLSTK
DO 310 J=1,ND
J0=(J-1)*ND
DO 310 I=1,ND
I0=J0+I
K0=0
K1=0
DO 300 K=MC1,ND2
K0=K0+1
IF(MK(I0,K).EQ.0)GO TO 300
K1=K1+1
MV(K1)=K0
MC(K1)=MK(I0,K)
300 CONTINUE
IF(K1.EQ.0)GO TO 305

```

```
CALL CRPHD1(KI,MV,MC,M4(I,J))
GO TO 310
305 M4(I,J)=0
310 CONTINUE
MR>IDET4(M4)
CALL PDCOD(MR,MV,MC,KT,100)
IF(KT.GE.99)GO TO 380
I0=0
DO 350 I=1,KT
350 I0=IGCD(I0,MC(I))
IF(I0.NE.1)GO TO 400
PRINT 2,YDEC,IAC,JDEC,JAC,KA
DO 360 I=1,KT
I0=0
J0=MV(I)
355 IF(J0.EQ.0)GO TO 360
K0=J0/24
L0=J0-24*K0
J0=K0
I0=I0+1
MP(I0)=L0
GO TO 355
360 PRINT 3,MC(I),(MP(J),J=1,I0)
GO TO 35
380 PRINT 4,IDEC,IAC,JDEC,JAC
GO TO 35
400 CONTINUE
GO TO 35
END
```

```

PROGRAM ACCHK
0 DIMENSION LRI(4,8),LRE(4,8),LOTW(7),LTI(4,7),LTE(4,7),MRI(5,25)
1,M1(5,5),MR(40,25),M2(5,5),MT(5,5,7),MRT(35,25),ISOC(2)
COMMON MA(5,5,5),ND,NDMU
1 FORMAT(13)
2 FORMAT(5I3,2A4)
3 FORMAT(5HISKIP,14,1X,5HACREC)
4 FORMAT(5HONEXT,13,35H ARITH CRYST CLASSES ISOMORPHIC TO ,2A4
171X,13,11H GENERATORS,10X,13,9H RELATORS,10X,13,13H TOR 1ST ELEM)
6 FORMAT(4H REL,13,10X,10I3)
7 FORMAT(8H ITD ORD,13,6X,10I3)
8 FORMAT(13H0ARITH CRYST CLASS,14,10X,3HDIM,13,10X,9H1SO CLASS,2A4)
9 FORMAT(1H /4RSGEN,13)
10 FORMAT(15X,5I3)
11 FORMAT(1H /9H1TORT ORD,13)
12 FORMAT(14H0TAPE SEQUENCE,14)
13 FORMAT(13H ***FOUL DATA, 2I3)
NLRAC=0
CALL MHPY1(5,5,5,0,0,0,1)
MNL=4
C SKIP PREVIOUS RECORDS
READ 1,NTAC
PRINT 3,NTAC
IF(NTAC)30,30,20
20 DO 25 I=1,NTAC
25 READ TAPE 1,I,IFL
C GET ISO CLASS
30 READ 2,IFL,KAC,MU,NU,NTW,ISOC
IF(IFL-999)300,31,300
31 PRINT 4, KAC, ISOC,MU,NU,NTW
DO 35 I=1,NU
READ 1,(LRI(J,I),LRE(J,I),J=1,MNL)
35 PRINT 6,I,(LRI(J,I),LRE(J,I),J=1,MNL)
DO 40 I=1,NTW
READ 1,(LOTW(I),(LTI(J,I),LTE(J,I),J=1,MNL)
40 PRINT 7,LOTW(I),(LTI(J,I),LTE(J,I),J=1,MNL)
50 DO 500 IACB1,KAC
NTAC=NTAC+1
READ 1,IFL,NAC,ND
IF(IFL-888)500,60,560
60 CALL MHPY2(ND,ND,ND,1)
PRINT 8,NAC,ND,ISOC
DO 65K=1,MU
PRINT 9,K
READ 1,((MA(I,J,K),J=1,ND),I=1,ND)
DO 65 I=1,ND
65 PRINT 10,(MA(I,J,K),J=1,ND)
NDMU=ND*MU
NDNU=ND*NU
NDNTW=ND*NTW
C CHECK RELATORS
DO 125 INU=1,NU
CALL WORDV(LRI(1,INU),LRE(1,INU),MRI,M1)
DO 120 I=1,ND
DO 120 J=1,ND

```

```

IF (I-J) 105, 110, 105
105 IF (M1(I,J)) 115, 120, 115
110 IF (M1(I,J)-1) 115, 120, 115
115 IFU=INU
IGU=1
GO TO 570
120 CONTINUE
I1=ND*(INU-1)
DO 125 I=1,ND
I2=I*I
DO 125 J=1,NDMU
125 MR(I2,J)=MRI(I,J)
C COMPUTE FOR 1ST ELEM AND CHECK ORDER
DO 280 ITW=1,NTW
I1=ND*(ITW-1)
CALL WORDV(LTE(1,ITW),LTE(1,ITW),MRI,MT(1,1,ITW))
CALL MPPZ(MT(1,1,ITW),MT(1,1,ITW),M1)
ISW=1
KI=LOTW(ITW)-1
GO TO 220
209 GO TO (210,215),ISW
210 CALL MPPZ(M1,M2)
ISW=2
GO TO 220
215 CALL MPPZ(M2,M3)
ISW=3
220 KI=KI-1
221 IF (KI) 225, 225, 209
225 GO TO (230,240),ISW
230 DO 235 I=1,ND
DO 235 J=1,ND
IF (I-J) 231, 232, 231
231 IF (M1(I,J)) 250, 235, 250
232 IF (M1(I,J)-1) 250, 235, 250
235 CONTINUE
GO TO 260
240 DO 245 I=1,ND
DO 245 J=1,ND
IF (I-J) 241, 242, 241
241 IF (M2(I,J)) 250, 245, 250
242 IF (M2(I,J)-1) 250, 245, 250
245 CONTINUE
GO TO 260
250 IFU=ITW
IGU=2
GO TO 570
260 PRINT I1,LOTW(ITW)
DO 270 I=1,ND
270 PRINT 10,(MT(I,J,ITW),J=1,ND)
DO 280 I=1,ND
I2=I*I
DO 280 J=1,NDMU
280 MR(I2,J)=MRI(I,J)
C OUTPUT
WRITE TAPE 1,NLRAC,ND,NAC,ISOC,MU,NU,NTW,NDMU,NDNU,NDNTW,((LRI(I,J

```

```
1),LRE(I,J),I=1,MWL),J=1,NU),(LOTW(I),(LTI(J,I),LTE(J,I),J=1,MWL),  
2I=1,NTW),((MA(I,J,K),I=1,ND),J=1,ND),K=1,HU),((MT(I,J,R),I=1,ND)  
3,J=1,ND),K=1,NTW),((MR(I,J),I=1,NDNU),J=1,NDMU),((MRT(I,J),I=1,  
4NDNTW),J=1,NDMU)  
500 PRINT 12,NTAC  
GO TO 30  
560 IFU=0  
IGU=0  
570 PRINT 13,IFU,IGU  
600 REWIND 1  
STOP  
END
```

```

PROGRAM HILB18
COMMON MR(40,25),MRI(40,25),MK(25,25),MCV(25),ME(25),MCC(25),
MFBS(25),MRFS(25),MFBST(35),MRT(35,25),LRI(4,8),LRE(4,8),LOTW(7),
2LTI(4,7),LTE(4,7),MA(5,5,5),MT(5,5,7),MCW(25),ISOC(2)
1 FORMAT(26I3)
20FORMAT(3H0AC,I4,5X,3HDIM,I4,5X,4HISOC,2A4/
13H H2,4X,3HORD,I5,4X,4HRANK,I5)
3 FORMAT(6H INVAR,25I4)
4 FORMAT(1H /9H*K-MATRIX)
5 FORMAT(10X,25I4)
MWL=4
ICHP=-1
CALL DIAGK1(40,25,MRI,KK,MK)
C SKIP RECORDS
READ I,I1,I2,NPTGP
101 DO 105 I=1,I1
105 READ TAPE 1
115 DO 120 I=1,I2
READ TAPE 2,NLRAC
116 DO 119 J=1,NLRAC
119 READ TAPE 2
120 CONTINUE
130 NPTGP=1
131 IF (IPTGP-NPTGP)200,200,140
140 REWIND 1
REWIND 2
STOP
200 READ TAPE 1 ,NLRAC,ND,NAC,ISOC,MU,NU,NTW,NDMU,NDNU,NDNTW,((LRI(I,
1),LRE(I,J),I=1,MWL),J=1,NU), (LOTW(I), (LTI(J,I),LTE(J,I),J=1,MWL),
2),I,NTW), ((MA(I,J,K),I=1,ND),J=1,ND),KSI,MU), ((MT(I,J,K),I=1,ND
3),J=1,ND),K=1,NTW), ((MR(I,J),I=1,NDNU),J=1,NDMU), ((MRT(I,J),I=1,
4NDNTW),J=1,NDMU)
DO 210 J=1,NDMU
DO 210 I=1,NDNU
110 MRI(I,J)=MR(I,J)
CALL DIAGK(NDNU,NDMU)
C GET S2
NOH2=1
KPK=0
220 HERO=MR1(KK,KK)
DO 240 IK=1,KK
K1=MR1(YK,IK)
NOH2=NOH2*K1
NOV(IK)=HERO/K1
ME(IK)=K1
IF (K1.EQ.1) GO TO 240
KPK=KPK+1
240 CONTINUE
PRINT 2,NAC,ND,ISOC,NOH2,KPK
245 PRINT 3,(ME(I),I=1,KK)
PRINT 4
DO 246 I=1,NDMU
246 PRINT 5,MR(I,J),J=1,NDMU)
260 NLRAC=NOH2-1
WRITE TAPE 2,NLRAC,ND,NAC,ISOC,MU,NU,NTW,NDMU,NDNU,NDNTW,((LRI(I,

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1) ,LRE(I,J),I=1,MWL),J=1,NDU), (LOTW(I), (LTI(J,I),LTE(J,I),J=1,MWL),
2I=1,NTW), ((MA(I,J,K),I=1,ND),J=1,ND),K=1,MU), ((TMT(I,J,R),I=1,ND)
3,J=1,ND),K=1,NTW),KK ,NOH2,MERO,(ME(I),I=1,KK), ((MK(I,J),I=1,NDMU)
4),J=1,NDMU)
IF(NLRAC) 400,400,300
300 DO 305 I=1,KK
305 MCC(I)=0
KEQ=1
310 KINC=1
315 MCI=MCC(KINC)*MCV(KINC)
IF(MCI-MERO)325,320,320
320 KINC=KINC+1
IF(KINC-KK)315,315,400
325 KIN=KINC-1
330 DO 335 I=1,KIN
335 MCC(I)=0
340 MCC(KINC)=MCI
CALL HNPY1(25,25,25,NDNU,KK ,1,0,MK,MCC,MFBS)
CALL HNPY1(35,25,25,NDNTW,NDMU,1,0,MRT,MFBS,MFBST)
CALL HNPY1(40,25,25,NDNU,NDMU,1,0,MR,MFBS,MRFS)
DO 360 I=1,NDNU
360 MRFS(I)=MRFS(I)/MERO
DO 385 I=1,KK
385 MCW(I)=MCC(I)/MCV(I)
WRITE TAPE 2,(ICHF,I=1,5),(MCW(I),I=1,KK),(MFBS(I),I=1,NDMU),
I(MRFS(I),I=1,NDNU),(MFBST(I),I=1,NDNTW)
GO TO 310
400 IPTGP=IPTGP+1
GO TO 131
END

```

```

PROGRAM FINDFLAT
COMMON ISOC(2), LOTW(7), MT(5,5,7), LTC(7), MCW(25), MFBST(5,7), M1(5,5)
1 M2(5,5), M3(5,5), KL(7), MTT(5,5,7), MTL(5,5,7), MID(5,7)
2 MRHS(5), MLHS(5), LHS(5)
1 FORMAT(26I3)
20 FORMAT(12H0 ARITH CLASS, I4, 5X, 3H0 I1, I2, 5X, 4HISOC, 2A4 /
110H TEST ELEM, I3, 10X, 7H02 RANK, I3, 5X, 5HORDER, I4)
3 FORMAT(10X, 11H0 COM CLASS, 25I3)
4 FORMAT(10X, 14H0 CRITICAL ELEM, I3, 10X, 5HPOWER, I3)
5 FORMAT(5X, 35H*****TORSION FREE*****ILRAC=, I4)
6 FORMAT(10X, 28H CRITICAL FOR ALL EXTENSIONS)
CALL MMPY1(5,5,5,0,0,0,1)
CALL DIAGL1(5,5,M2, KK, M3)
MWL=4
C SKIP RECORDS
READ 1, NSK, NDO
DO 20 I=1, NSK
READ TAPE 2, NLRAC
DO 20 J=1, NLRAC
20 READ TAPE 2
C NEXT ARITH CLASS
100 IF (NDO) 1000, 1000, 105
1050 READ TAPE 2, NLRAC, ND, NAC, ISOC, MU, NU, NTW, NDMU, NDN, NDNTW, ((KBG, KBG,
I1=1, NDL), J=1, NU), (LOTW(I), (KBG, KBG, I=1, MWL), J=1, NTW), (((KBG, I=1, ND
2), J=1, ND), K=1, ND), ((MT(I, J, K), I=1, ND), J=1, ND), K=1, NTW), KKK, NOH2,
3MERS
PRINT 2, NAC, ND, ISOC, NTW, KKK, NOH2
1 INLRAC) 106, 106, 110
106 NDO=NDO-1
GO TO 100
110 CALL DIAGL2(ND, ND)
DO 115 I=1, NTW
115 LTC(I)=0
DO 500 I=1, NLRAC
O READ TAPE 2, (KBG, I=1, 5), (MCW(I), I=1, KKK), (KBG, I=1, NDMU), (KBG, I=1,
INDN), ((MFBST(I, J), I=1, ND), J=1, NTW)
PRINT 3, (MCW(I), I=1, KKK)
DO 400 ITW=1, NTW
IF (LTC(ITW)) 120, 120, 130
120 LOR=1
LO=LOTW(ITW)
CALL MMPY2(ND, ND, ND, 1)
CALL MAT1(ND, LO-1, MT(1, 1, ITW), MTL(1, 1, ITW))
CALL MATP(ND, 1, MTT(1, 1, ITW), M2)
CALL DIAGL
IF (KK) 390, 390, 130
130 KL(ITW)=KK
CALL MMPY(M3, MTT(1, 1, ITW), MTL(1, 1, ITW))
DO 135 I=1, KK
135 MID(I, ITW)=MRHS(I)=M2(I, I)
LTC(ITW)=1
GO TO 160
160 LO=LOTW(ITW)
KK=KL(ITW)
DO 165 I=1, KK

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```

165 MRHS(I)=MTD(I,ITW)
    LOR=1
166 CALL MNPY2(ND,ND,1,0,MTTL(1,1,ITW),MFBST(1,ITW),MLHS)
    DO 170 I=1,KK
170 MLHS(I)=MLHS(I)/MERO
200 L=1
    DO 215 I=1,KK
        KZ=MLHS(I)
        IF (KZ) 205,215,205
205 K2=MRHS(I)/IGCD(MRHS(I),KZ)
        K3=LO-(LO/K2)*K2
        IF (K3) 400,210,400
210 L=(L*K2)/IGCD(L,K2)
215 CONTINUE
        IF (L-LO) 220,400,400
220 IF (L-1) 300,380,225
225 IF (LOR-1) 230,230,240
230 CALL MATP(ND,1,MT(1,1,ITW),MI)
    CALL MNPY2(ND,ND,1,0,MTT(1,1,ITW),MFBST(1,ITW),LHS)
    DO 232 I=1,ND
232 LHS(I)=LHS(I)/MERO
240 LOR=LOR*L
    LO=LO/L
    CALL MNPY2(ND,ND,ND,1)
    CALL MATP(ND,LV1,K2)
    CALL MATP(ND,1,N2,MI)
    CALL MATP(ND,LO-1,MI,N2)
    CALL D11GL
        IF (KK) 390,390,250
250 DO 255 I=1,KK
255 MRHS(I)=M2(I,I)
    CALL MNPY2(ND,ND,1,0,M3,LHS,MLHS)
    GO TO 200
380 PRINT 4,ITW,LOR
    GO TO 500
390 PRINT 4,ITW,LOR
    PRINT 6
        K2=1000
    DO 395 I=K2,NLRAC
395 READ TAPE 2
    GO TO 106
400 CONTINUE
    PRINT 5 ICC
500 CONTINUE
    GO TO 106
1000 REWIND 2
    STOP
    END

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```

PROGRAM ABELR
COMMON LRI(4,8),LRE(4,8),MA(5,5,5),MRS(10,33),MR(10,33),ME(25),
IMCC(25),MRFS(40),ISOC(2)
INTEGER GSG
1 FORMAT(26I3)
2 FORMAT(1H0,10X,3HACC,14,10X,3HDIM,14,10X,2HHZ,5X,5HORDER,14,5X,
14HRANK,14,10X,10HINVARIANTS,2X,25I3)
3 FORMAT(12H0COHOM CLASS,69X,22HEXTENSION MADE ABELIAN)
4 FORMAT(1H*,25I3)
5 FORMAT(81X,10I3)
MWL=4
CALL DERGE(10,MR,KK)
C SKIP RECORDS
READ 1,NSK,ND0
DO 10 I=1,NSK
READ TAPE 2,NLRAC
DO 10 J=1,NLRAC
10 READ TAPE 2
15 IF(ND0)16,16,20
16 REWIND 2
STOP
20 READ TAPE 2,NLRAC,ND,NAC,ISOC,MU,NU,NTW,NDMU,NDNU,NDNTW,((LRI(I,J),
1,LRE(I,J),I=1,MWL),J=1,NU), (GSG,(GSG,GSG,I=1,MWL),J=1,NTW),
2(((MA(I,J,K),I=1,ND),J=1,ND),K=1,MU),(((GSG,I=1,ND),J=1,ND),K=1,
SNTW),KKR,NCHZ,GSG,(ME(I),I=1,KRK)
PRINT 2,NAC,ND,NCHZ,KKR,(ME(I),I=1,KRK)
PRINT 3
C SET UP CONSTANT PART OF MR
NR1=ND+1
NR2=ND+NU
NC1=NU+1
NC2=NU+ND*MU
J1=NU-ND
DO 120 J=1,MU
J1=J1+ND
DO 130 J2=1,ND
J3=J1+J2
DO 129 I=1,ND
129 MRS(I,J3)=MA(I,J2,J)
130 MRS(J2,J3)=MRS(J2,J3)-1
DO 140 I=NR1,NR2
DO 140 J=1,NC2
140 MRS(I,J)=0
DO 165 J=1,NU
DO 160 J1=1,MWL
J2=LRI(J1,J)
IF (J2)165,165,150
150 J3=ND+J2
160 MRS(J3,J)=MRS(J3,J)+LRE(J1,J)
165 CONTINUE
KERR1
C SPLIT EXTENSION
DO 210 I=1,KRK
210 MCC(I)=0
DO 215 J=1,NU

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```

      DO 215 I=1,ND
215  MR(I,J)=0
      GO TO 230
C NON TRIVIAL EXTENSIONS
2250 READ TAPE 2, (ICHP, I=1,5), (MCC(I), I=1, KKK), (G8G, I=1, NDMU),
      I, MRFS(I), I=1, NDMU)
      II=0
      DO 245 J=1, NU
      DO 245 I=1, ND
      II=II+1
245  MR(I,J)=-MRFS(II)
230  DO 235 J=NC1, NC2
      DO 235 I=1, NR2
235  MR(I,J)= MRS(I,J)
      DO 240 I=NR1, NR2
      DO 240 J=1, NU
240  MR(I,J)= MRS(I,J)
260  CALL DIAG(NR2, NC2)
      DO 265 I=1, NR2
      IF (MR(I,I)-1) 264, 265, 266
264  II=I
      GO TO 266
265  CONTINUE
266  PRINT 4, (MCC(I), I=1, KKK)
      PRINT 5, (MR(I,I), I=1, NR2)
      IF (KEQ-NOH2) 270, 275, 275
270  KEQ=KEQ+1
      GO TO 225
275  NDC=NDC-1
      GO TO 15
      END

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PROGRAM ISG
COMMON MCOM(2700),MLB(15,15),IB(10,3),MA(5,5,3),MRI(15,15),
1 MK(25,25),MFBS(15),NFBS(15),ND,NDMU,MPE(25)
2,MEH(15),MKH(15,15)
0DIMENSION MX(75,25),MP(15,45),MKB(25,45),MLC(25,25),MLG(25),
1 MKB1(45,45),MB(5,5,3),LAI(4),LAE(4)
2 DIMENSION MED(25,25)
0EQUIVALENCE (MCOM,MK,MP,MLC,LAI),(MCOM(5),LAE),(MCOM(1876),MB),
1 (MCOM(676),MKB,MKB1),(MCOM(1801),MLD),(MCOM(626),MLG)
1 FORMAT(25I3)
2 FORMAT(4H0ACC,4I4)
3 FORMAT(5X,5H$IMCL,10I4)
4 FORMAT(1X/1H*,10X,4HAUT ,I4)
5 FORMAT(25X,0I4)
6 FORMAT(1X,2I4,1X,12HNO PART SOLN )
7 FORMAT(5H ERRORS,15)
MAXQ=10
MUL=4
NZO=0
C GET HOLONOMY
25 READ 1,NAC,NAUT,NEQ,KORD
PRINT 2,NAC,NAUT,NEQ,KORD
IF(NAC,GE,1)GO TO 30
26 REWIND 10
WRITE TAPE 11,(NZO,I=1,100)
REWIND 11
STOP
30 READ TAPE 10,NLRAC,ND,KAC,IO,IO,MU,IO,II,NDMU,KO,KO,((KO,KO,I=1,
1 MUL),J=1,IO),(IO,(IO,IO,I=1,MUL),J=1,II),(((MA(I,J,K),I=1,ND),J=1
2,ND),K=1,MU),(((IO,I=1,ND),J=1,ND),K=1,II),KH,IO,MERO
3,(MEH(I),I=1,KH),((MKH(I,J),I=1,NDMU),J=1,NDMU)
IF(KAC-NAC)31,40,35
31 DO 32 I=1,NLRAC
32 READ TAPE 10
GO TO 30
35 MER =1
GO TO 950
C GET SIM CLASSES
40 DO 45 I=1,NEQ
READ 1,(IB(J,I),J=1,MAXQ)
45 PRINT 3,(IB(J,I),J=1,MAXQ)
KORDERO=KORD*MERO
ND2=ND**2
ND2MU=ND2*MU
IO=1
DO 445 J=1,KH
JI=MEH(J)
MEH(J)=IO
=IO*JI
JI=KORDERO/JI
DO 445 I=1,NDMU
445 MKH(I,J)=MKH(I,J)*JI
KSI=KSI+1
KOUT=0
C CENTRALIZER--IDENTITY AUTOMORPHISM

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      CALL NITTY(ND,ND,MU,MA,MA,MX)
      GO TO 60
C NEXT AUTOMORPHISM
  50 KAUT=KAUT+1
      PRINT 4,KAUT
C SET UP MRI
      CALL MXPY1(S,D,S,ND,ND,ND,1)
      KI=1
      DO 55 K=1,MU
      READ 1,(LAI(I),LAE(I),I=1,MWL)
      PRINT 5,(LAI(I),LAE(I),I=1,MWL)
      CALL WORDV(LAI,LAE,MRI(KI),KB(1,1,K))
  55 KI=KI+ND
      CALL NITTY(ND,ND,MU,MA,MB,MX)
  60 CALL DIAR(25,MK,75,MX,KKS,ND2MU,ND2)
      NVI=KKS+1
      NV=ND2-KKS
C NEXT SIMCL
      KEQ=1
  99 IEQ=1
 100 JEQ=IEQ+1
      JLRAC=IS(JEQ,KEQ)
      IF(JLRAC.EQ.0)GO TO 355
      YLRAC=KLRAC-IS(IEQ,KEQ)
      IS=1
      GO TO 450
 120 IF(KAUT.GT.0)GO TO 130
      DO 125 I=1,NDMU
 125 MFBS(I)=MFBS(I)
      GO TO 550
 130 CALL MXPY2(NDMU,NDMU,1,0,MRI,MFBS,MFBS)
      GO TO 550
C SET UP PS
 140 K0=0
      K3=NV+ND
      K4=K3+1
      K5=K3+NDMU
      DO 180 K=1,MU
      DO 175 I=1,ND
      K1=K0+I
      DO 160 L=1,NV
      K2=KKS+L
      AL=0
      J1=1
      DO 145 J=1,ND
      J2=K0+J
      ML=K1+MK(J1,K2)*MFBS(J2)
 145 J1=J1+ND
 150 MP(K1,L)=ML
      DO 150 J=1,ND
      L=J+NV
 160 MP(K1,L)=MA(I,J,K)
      L=I+NV
      MP(K1,L)=MP(K1,L)-1
      DO 165 L=K4,K5

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165 MP(K1,L)=0
    L=K3+K1
    MP(K1,L)=KORDERO
175 CONTINUE
180 K0=K0+ND
    CALL DIAKL(15,MLB,45,MKB1,15,MP,KKB,NDMU,K5)
    DO 185 I=1,KKB
185 MPE(I)=MP(I,I)
    KKB1=KKB+1
    DO 182 J=1,K5
    DO 182 I=1,NV
182 MKB(I,J)=MKB1(I,J)
    CALL DIALI(25,MLC,25,MKS(1,KKB1),KKC,NV,K5-KKB)
    CALL MMPY1(25,25,25,NV,KKC,KKC,0,MLC,MKB(1,KKB1),MLC)
    CALL MMPY2(ND2,NV,KKC,0,MK(1,NV1),MLC,MLD)
    CALL MMPY2(ND2,NV,KKB,0,MK(1,NV1),MKB,MLC)
205 KLRAC=JLRAC
    I0=2
    GO TO 450
220 CALL MMPY2(NDMU,NDMU,1,0,MLB,MFBS,NFBS)
    GO TO 550
300 DO 310 I=1,KKB
    K1=NFBS(I)
    K2=K1/MPE(I)
    K1=K1-K2*MPE(I)
    IF(K1.NE.0)GO TO 340
310 NFBS(I)=K2
    DO 320 I=KKB1,NDMU
    IF(NFBS(I).NE.0)GO TO 340
320 CONTINUE
    CALL MMPY1(25,25,25,ND2,KKB,1,0,MLC,NFBS,MLG)
    WRITE TAPE 11,NAC,KAUT,ILRAC,JLRAC,ND,KKC,((MLD(I,J),I=1,ND2),
    LOB,KKC),(MLG(I),I=1,ND2)
    GO TO 350
340 PRINT 6,ILRAC,JLRAC
350 IF(I0.GE.MAXQ)GO TO 351
    JEQ=JEQ+1
    KLRAC=IB(JEQ,KEQ)
    IF(ILRAC.EQ.0)GO TO 351
    GO TO 205
351 JEQ=JEQ+1
    IF(JEQ.GE.MAXQ)GO TO 355
    GO TO 100
355 KEQ=KEQ+1
    IF(KEQ.LE.NEQ)GO TO 99
    IF(KAUT.LT.NAUT)GO TO 50
    GO 360 I=1,NLRAC
360 READ TAPE 10
    GO TO 25
450 I=KKB1
    DO 455 J=1,KH
    J0=J-1
    J1=KKB(I,J)
    KLRAC(I)=K1+KLRAC/J1
455 KLRAC=KLRAC-K1*J1

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```
CALL NMPY1(15,15,15,NDMU,KH,1,0,NKH,NFBS,MPBS)
GO TO(120,220),10
550 DO 556 I=1,NDMU
    II=NFBS(I)
552 IF(II)554,556,556
554 II=II*KORDERO
    GO TO 552
556 NFBS(I)=II-KORDERO*(II/KORDERO)
    GO TO(140,300),10
950 PRINT 7,NER
    GO TO 26
END
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PROGRAM ISH
COMMON MSTK(1005),MLD(5,5,25),MLG(5,5),MC(100),MV(100),
IMX(4,4),NB(25)
EQUIVALENCE (KAUT,NAUT)
1 FORMAT(4HGRAC,3I4)
2 FORMAT(9H VAR OVFL,3I4)
3 FORMAT(9H POSS ISO/IZK,5HCOEFS,5X,9HVARIABLES)
4 FORMAT(5X,1I0,5X,25I0)
5 FORMAT(8H NON ISO,2I4)
LAUT=0
LILR=0
LAC=0
10 READ TAPE 11,NAC,KAUT,ILRAC,JLRAC,ND,KKC,((MLD(I,J,K),I=1,ND),J=
1,ND),K=1,KKC), (MLG(I,J),I=1,ND),J=1,ND)
IF(NAC.GT.0)GO TO 20
REWIND 11
STOP
20 PRINT 1,NAC,NAUT,ILRAC,JLRAC,KKC
IF(NAC.EQ.0.AND.LAC.AND.ILRAC.EQ.0.LILR.AND.NAUT.EQ.LAUT)GO TO 60
CALL CLSTK
DO 50 I=1,ND
DO 50 J=1,ND
KI=0
DO 40 K=1,KKC
K2=MLD(I,J,K)
IF(K2.EQ.0)GO TO 40
KI=KI+1
MV(KI)=K
MC(KI)=K2
40 CONTINUE
IF(KI.LE.23)GO TO 45
PRINT 2,KI,I,J
GO TO 10
45 K2=MLG(I,J)
IF(K2.EQ.0)GO TO 47
KI=KI+1
MV(KI)=0
MC(KI)=K2
47 IF(KI.GT.0)GO TO 49
MX(I,J)=0
GO TO 50
49 CALL CRPHD1(KI,MV,MC,MX(I,J))
50 CONTINUE
GO TO 100
60 DO 80 I=1,ND
DO 80 J=1,ND
IF(MX(I,J).EQ.0)GO TO 75
CALL HOMOG(MX(I,J))
75 K2=MLG(I,J)
IF(K2.EQ.0)GO TO 80
CALL UNHOG(MX(I,J),K2)
80 CONTINUE
100 NP=IDETA(MX)
CALL PDCGD(NP,MV,MC,KT,100)
CALL FRIPOL(NP)

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IF(KT.LT.99)GO TO 105
PRINT 2,KT
GO TO 200
105 IO=0
K4=0
DO 110 I=1,KT
IF(MV(I).EQ.0)GO TO 109
IO=IGCD(IO,MC(I))
GO TO 110
109 K4=MC(I)
110 CONTINUE
IF(IGCD(IO,K4).NE.1)GO TO 140
K2=K4+1
K1=K4-1
IF(IO.NE.0)GO TO 120
IF(K1.EQ.0.OR.K2.EQ.0)GO TO 125
GO TO 140
120 IF((K1/IO)*IO.NE.0.AND.K2-(K2/IO)*IO.NE.0)GO TO 140
125 PRINT 3
DO 160 I=1,K1
IO=0
JO=MV(I)
155 IF(JO.EQ.0)GO TO 160
K0=JO/24
L0=JO-24*K0
JO=K0
IO=IO+1
MV(IO)=L0
GO TO 155
160 PRINT 4,MC(I),(MV(J),J=1,IO)
GO TO 200
140 PRINT 5,IO,K4
200 LAST=NAUT
ILR=ILRAC
LAC=NAC
GO TO 10
END
FINIS

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